

Facilitating Non-Intrusive In-Vivo Firmware Testing with Stateless Instrumentation

NDSS Symposium 2024

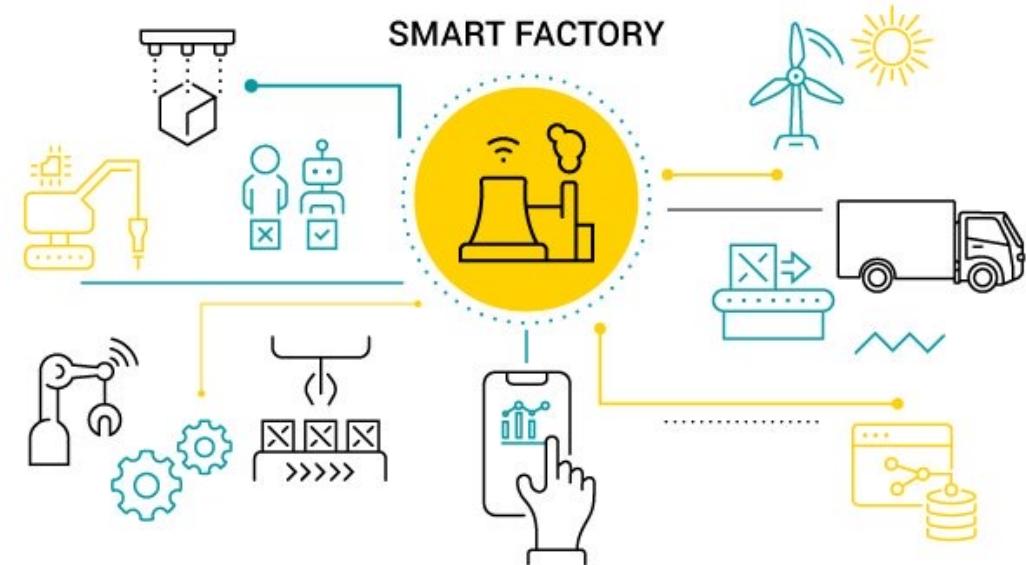
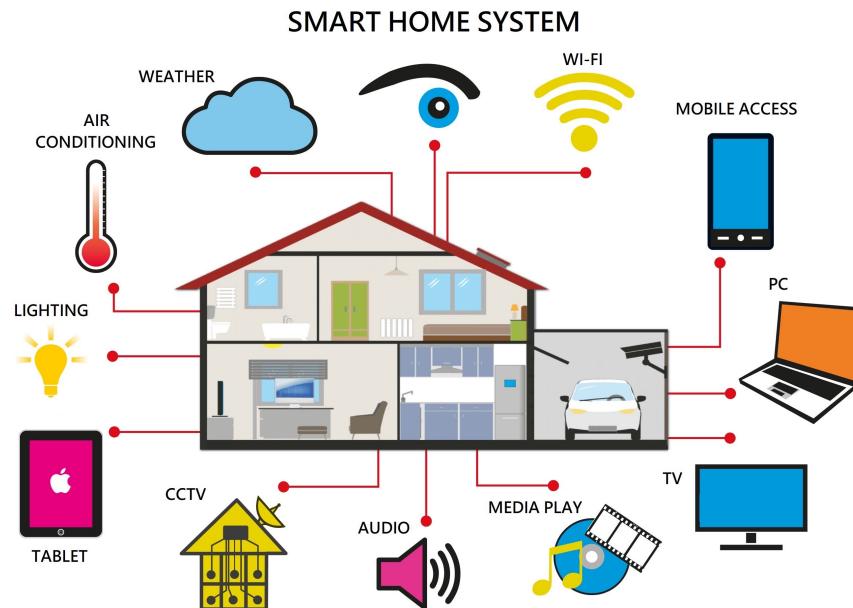
Jiameng Shi*, Wenqiang Li[#], Wenwen Wang*, Le Guan*

* The University of Georgia

Independent Researcher



MCU-based Devices are Ubiquitous



Img src: <http://visioforce.com/smarthome.html>

<https://www.avsystem.com/blog/smart-factory/>

But at the Same Time, They Are Extremely Vulnerable

Multiple Vulnerabilities in Treck TCP/IP Stack Could Allow for Arbitrary Code Execution

MS-ISAC ADVISORY NUMBER:

2020-171

DATE(S) ISSUED:

12/21/2020

OVERVIEW:

Multiple vulnerabilities have been discovered in Treck TCP/IP Stack, the most severe of which could result in arbitrary code execution. Treck TCP/IP Stack are networking protocols libraries specifically designed for **embedded systems** and are widely used. Successful exploitation of the most severe of these vulnerabilities could allow an attacker to execute arbitrary code in the context of the application. Depending on the privileges associated with the application, an attacker could install programs; view, change, or delete data; or create new accounts with full user rights. If this application has been configured to have fewer user rights on the system, exploitation of the most severe of these vulnerabilities could have less impact than if it was configured with administrative rights.

THREAT INTELLIGENCE:

There are currently no reports of these vulnerabilities being exploited in the wild.

But at the Same Time, They Are Extremely Vulnerable

Multiple Vulnerabilities in Treck

TCP/IP Stack

Arbitrary C

MS-ISAC ADVISORY NUMBER

2020-171

DATE(S) ISSUED:

12/21/2020

OVERVIEW:

Multiple vulnerabilities have been discovered in the Treck TCP/IP Stack, which is used for code execution. Treck TCP/IP Stack is a popular TCP/IP stack for industrial control systems and is widely used. Successful exploitation of these vulnerabilities could allow an attacker to execute arbitrary code in the context of the process running the stack. This could enable an attacker to install programs; view, change and control system configuration files; and change settings. If the stack has been configured to have fewer users, the impact of a successful exploit would have less impact than if it was configured to have more users.

THREAT INTELLIGENCE:

There are currently no reports of th



The hack attack led to failures in plant equipment and forced the fast shut down of a furnace

But at the Same Time, They Are Extremely Vulnerable

Multiple Vulnerabilities in Treck

TCP/IP Stack Arbitrary C

MS-ISAC ADVISORY NUMBER

2020-171

DATE(S) ISSUED:
12/21/2020

OVERVIEW:

Multiple vulnerabilities have been discovered in the Treck TCP/IP Stack, which is used for code execution. Treck TCP/IP Stack is a popular stack for industrial control systems and is widely used. Successful exploitation of these vulnerabilities could allow an attacker to execute arbitrary code in the context of the process running the stack. This could enable an attacker to install programs; view, change and control system data; and add, delete or update configuration files. If the stack has been configured to have fewer users and less memory protection, the exploit could have less impact than if it was configured to have more users and more memory protection.

THREAT INTELLIGENCE:

There are currently no reports of th



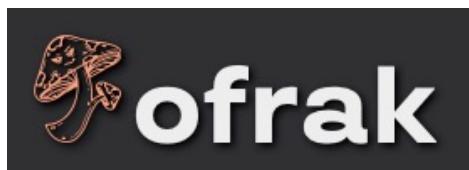
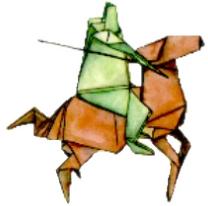
The hack attack led to failures in plant equipment and forced the fast shut down of a furnace.

Hacking risk leads to recall of 500,000 pacemakers due to patient death fears

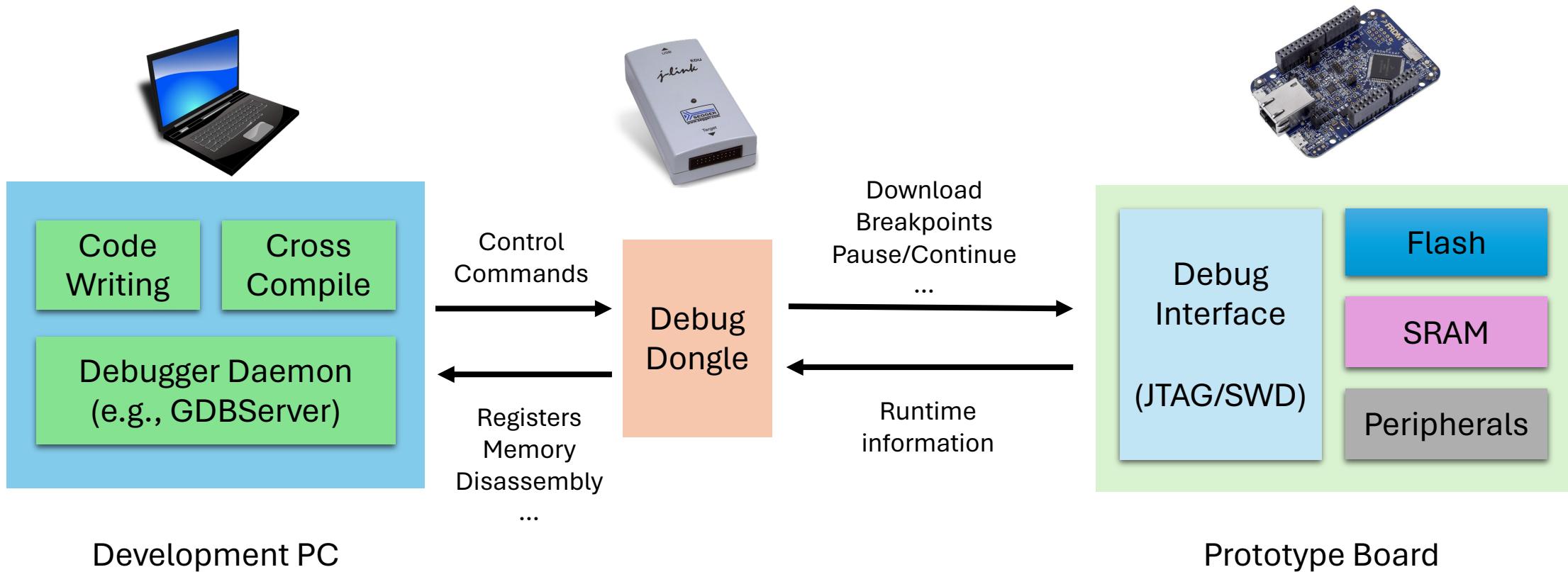
FDA overseeing crucial firmware update in US to patch security holes and prevent hijacking of pacemakers implanted in half a million people



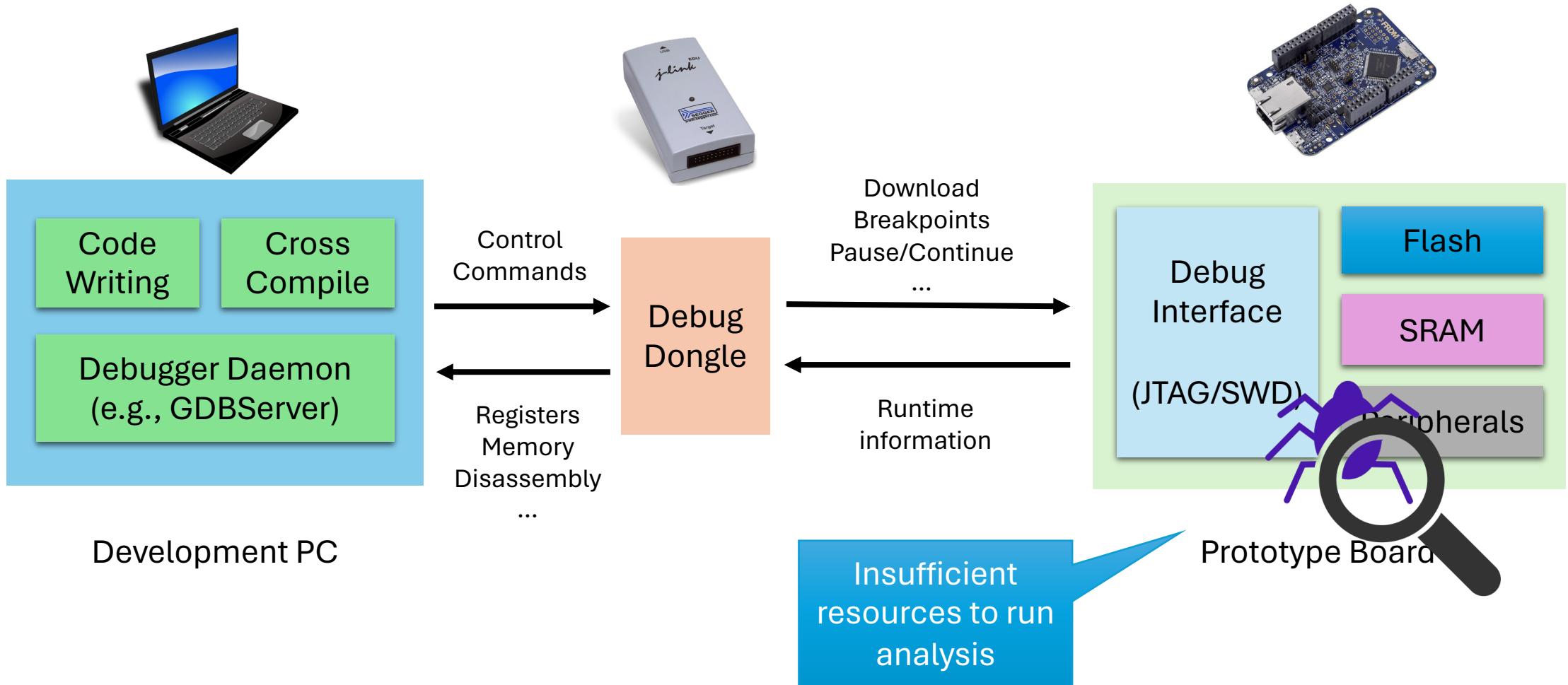
How to Effectively Find Firmware Bugs?



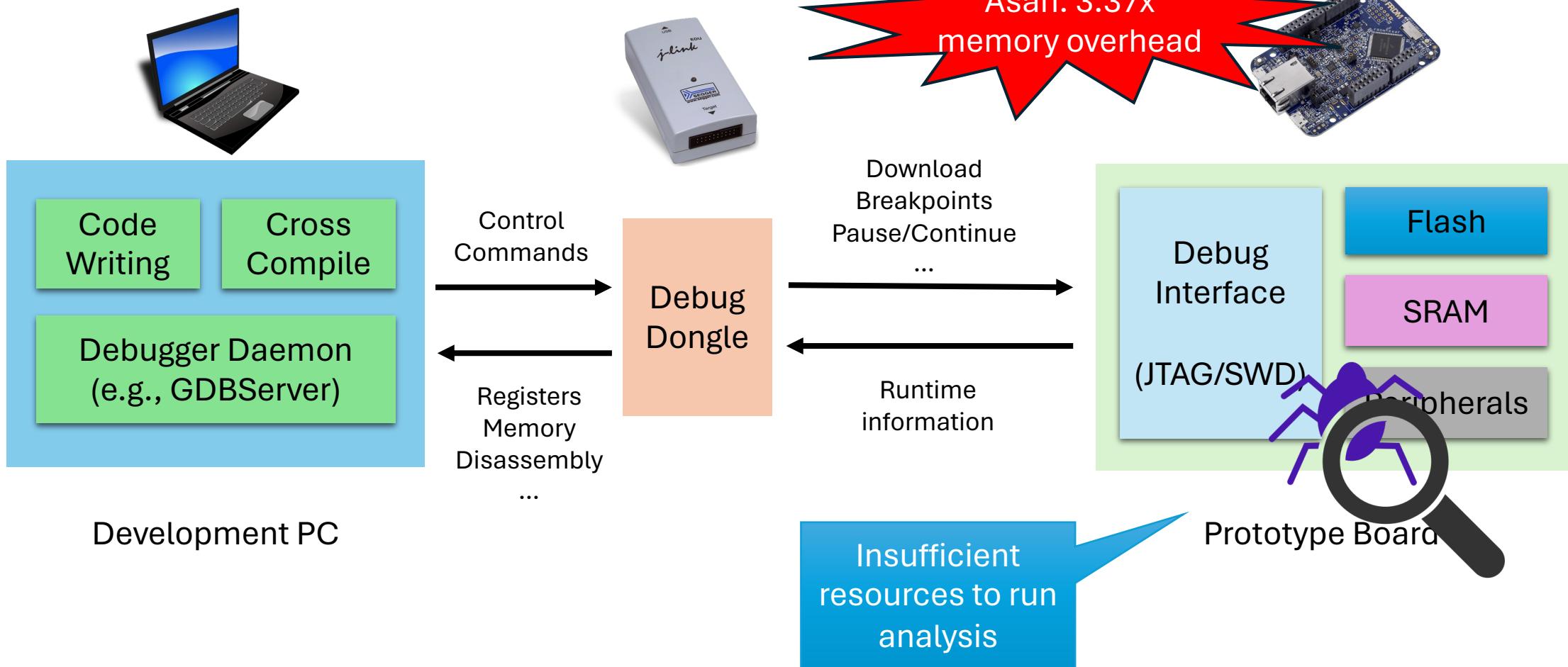
Firmware Development



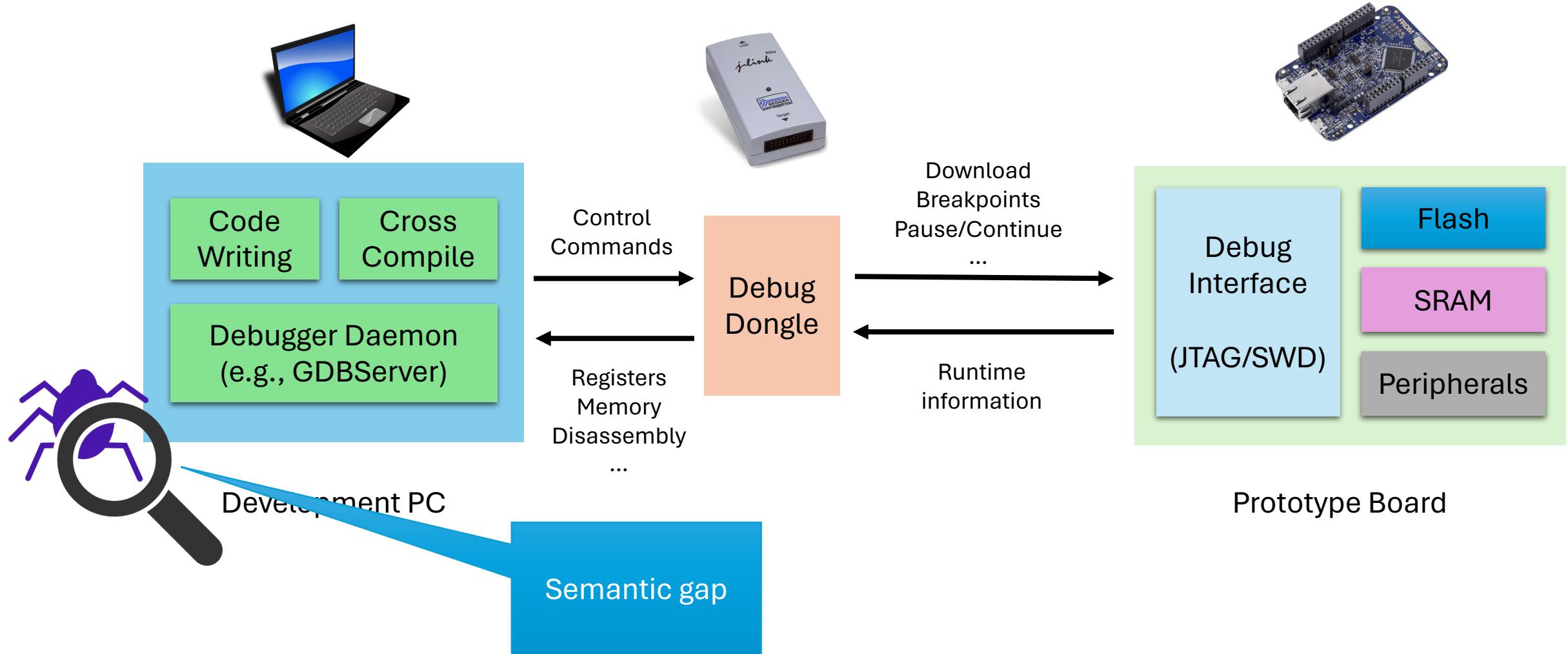
Analysis on the MCU Board?



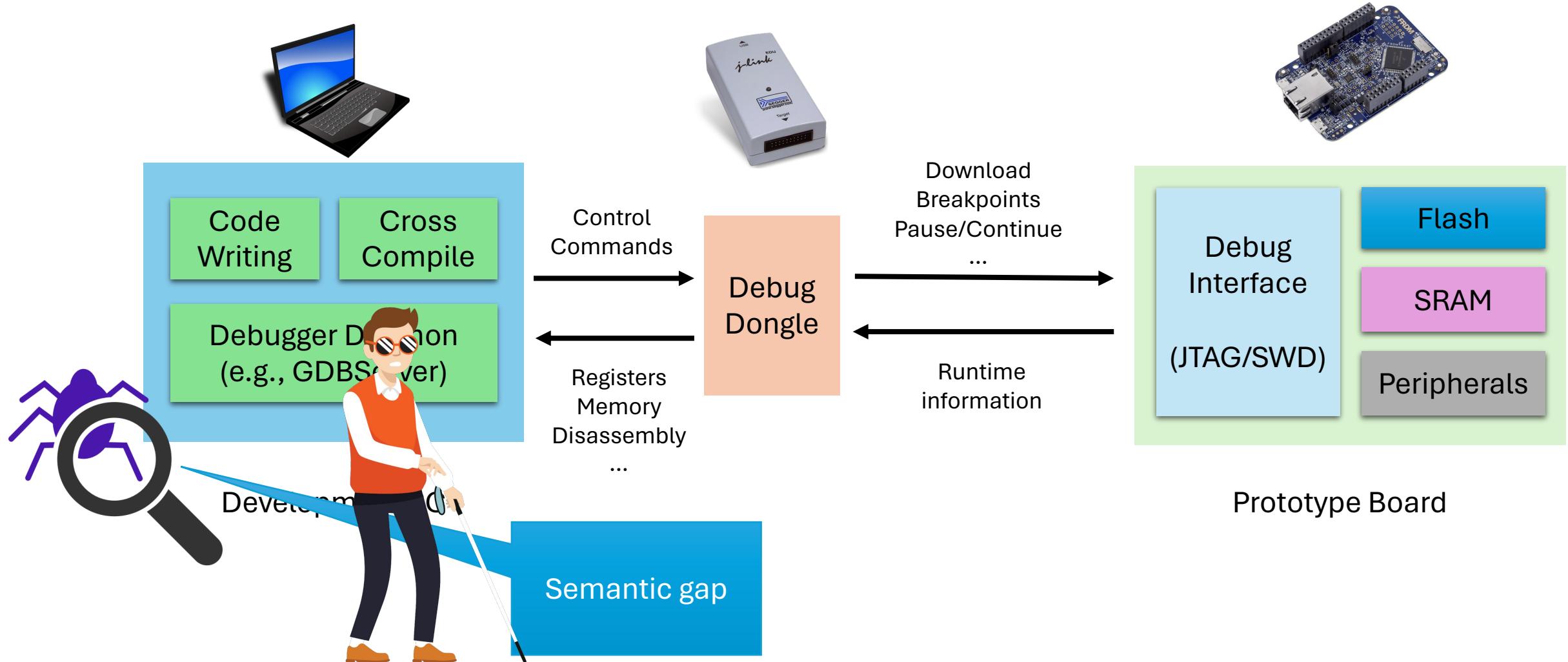
Analysis on the MCU Board?



Analysis on Development Workstation?

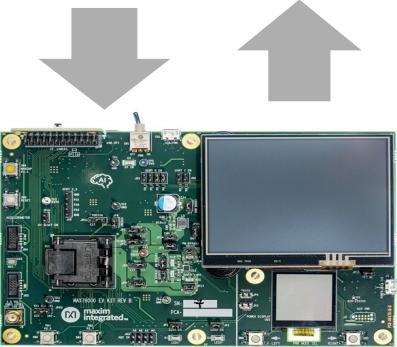


Analysis on Development Workstation?



State-of-the-Art Solutions

```
american fuzzy lop ++4.07a {default} ./opencv_test_wechat_qrcode [fast]
process timing
    run time : 0 days, 0 hrs, 1 min, 21 sec
    last new find : 0 days, 0 hrs, 0 min, 7 sec
last saved crash : none seen yet
last saved hang : none seen yet
cycle progress
    now processing : 23.1 (57.5%)
    runs timed out : 0 (0.00%)
stage progress
    now trying : splice 8
    stage execs : 10/12 (83.33%)
    total execs : 905
    exec speed : 7.89/sec (zzzz...)
fuzzing strategy yields
    bit flips : disabled (default, enable with -D)
    byte flips : disabled (default, enable with -D)
    arithmetics : disabled (default, enable with -D)
    known ints : disabled (default, enable with -D)
    dictionary : n/a
    havoc/splice : 3/108, 1/276
py/custom/rq : unused, unused, unused, unused
    trim/eff : 0.00%/237, disabled
map coverage
    map density : 8.45% / 15.91%
    count coverage : 3.36 bits/tuple
findings in depth
    favored items : 21 (52.50%)
    new edges on : 26 (65.00%)
    total crashes : 0 (0 saved)
    total tmounts : 0 (0 saved)
item geometry
    levels : 2
    pending : 38
    pend fav : 20
    own finds : 4
    imported : 0
    stability : 100.00%
[cpu000: 2x]
```



- μAFL [ICSE'22]
- GDBFuzz [ISSTA'23]
- Idea:
 - Run the firmware on device
 - Use **special** hardware to collect execution information
 - Stream the collected info on PC
- Limitations:
 - Limited to fuzzing
 - Only collect code coverage
 - Require hardware features
 - ETM, breakpoint

IPEA: In-vivo Probe, Ex-vivo Analysis

Non-intrusive

- IPEA seamlessly integrates into existing firmware development workflow and does not rely on additional hardware

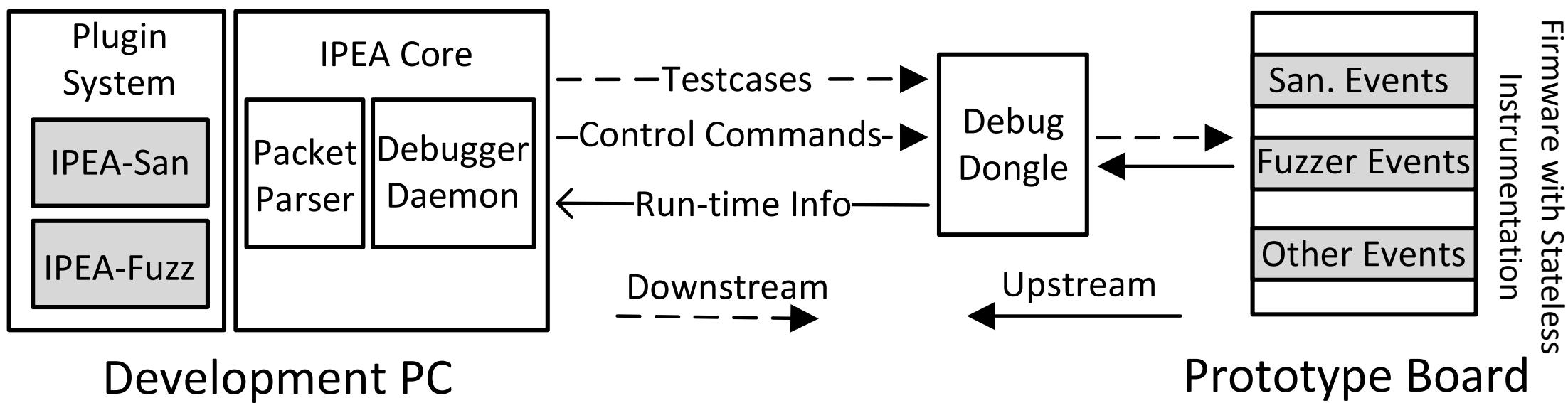
In-vivo

- The instrumented firmware collects arbitrary run-time information essential for firmware analysis, without dependent on special hardware

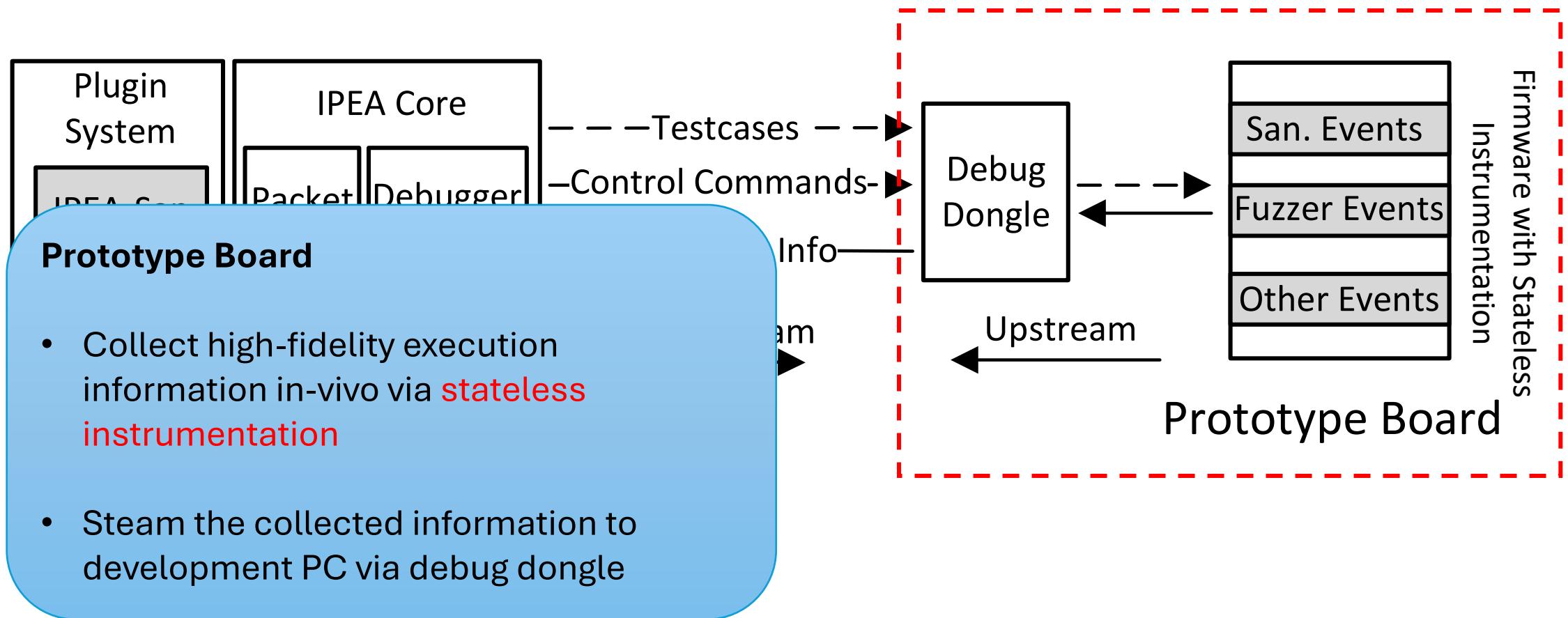
Lightweight

- Offload analysis to the PC, leaving the instrumentation on the device lightweight

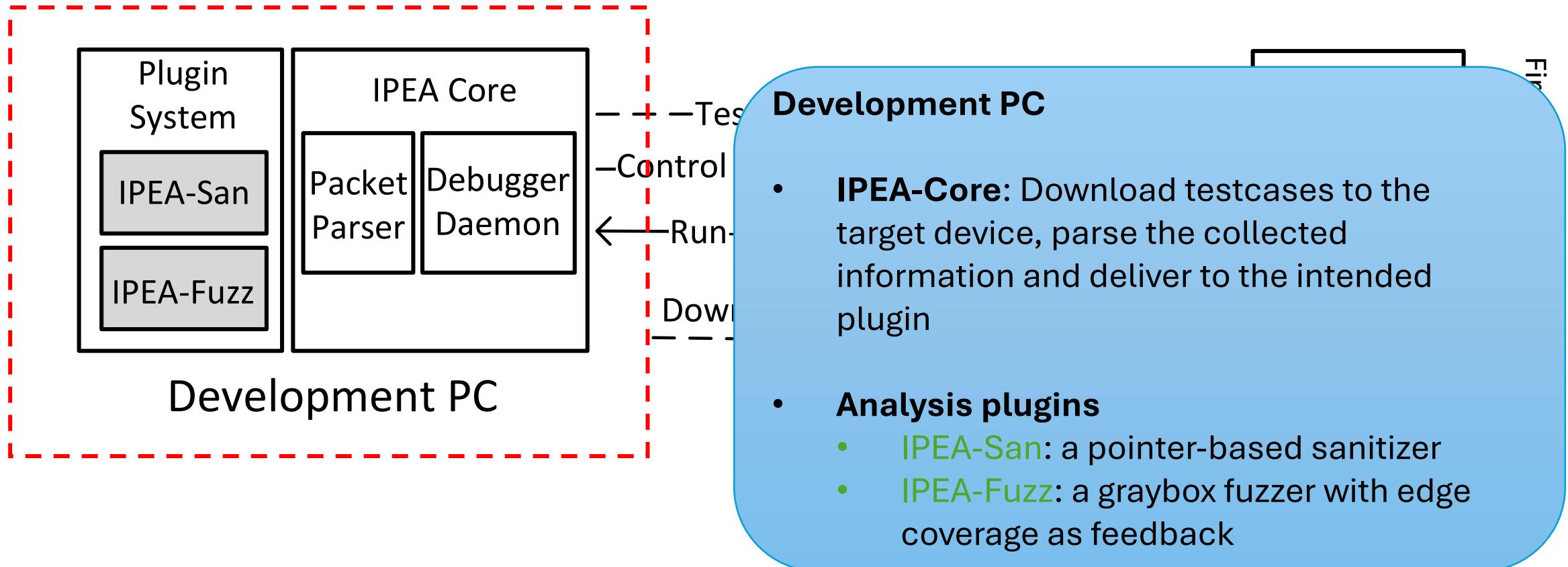
Design



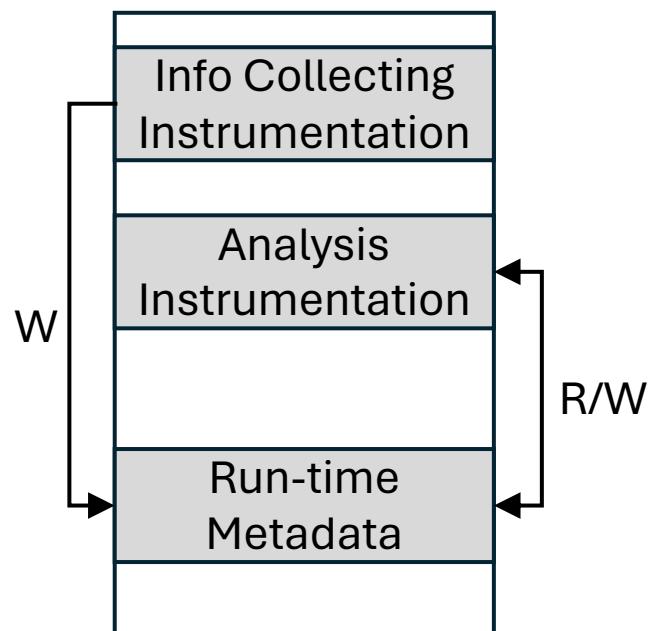
Design



Design

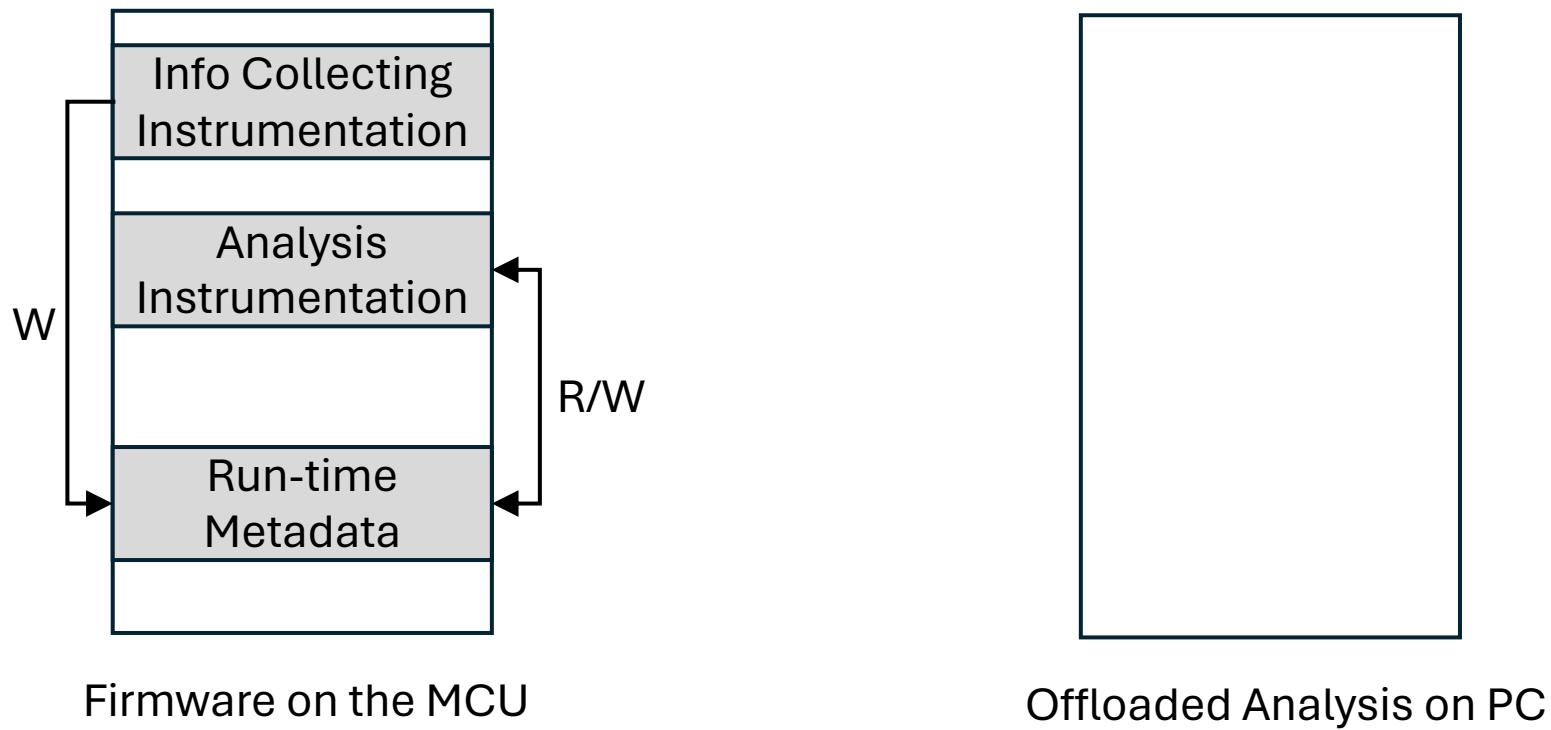


Stateless Instrumentation

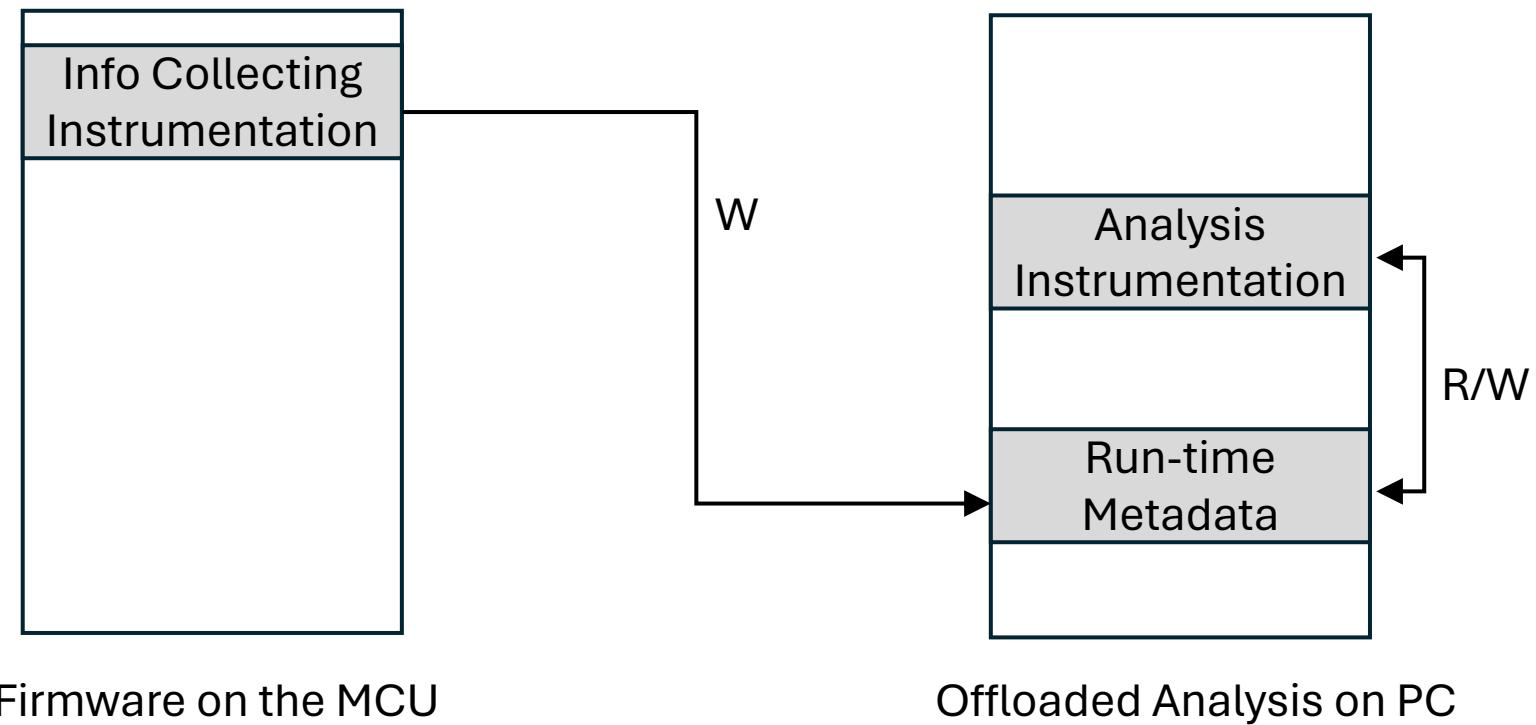


Firmware on the MCU

Stateless Instrumentation



Stateless Instrumentation



IPEA-San Plugin

- Pointer-based sanitizer
 - Avoid shadow memory
 - Each pointer is associated with metadata of its capability (bounds and validity)
 - E.g., SoftBound [PLDI'09]
- Memory tagging
 - A variant of pointer- based sanitizer
 - When an object is allocated, its memory and receiving pointer are given the same tag
 - All accesses to that memory must be made by a pointer having the same tag
 - E.g., Arm MTE
- We virtually extend the capability of MCUs by emulating an enhanced Arm MTE
 - Cooperation between the instrumented firmware and IPEA-San plugin on PC
 - Support unlimited number of tags compared with Arm MTE

Running Example – w/o IPEA-San

1. `ptr = malloc(size); // object creation`
- 2.
- 3.
4. `newptr = ptr + index; // pointer propagation`
- 5.
- 6.
- 7.
- 8.
9. `value = *newptr; // pointer dereference`

Running Example – w/o IPEA-San

1. `ptr = malloc(size); // object creation`
2. `ptr_base = ptr;`
3. `ptr_bound = ptr + size;`
4. `newptr = ptr + index; // pointer propagation`
- 5.
- 6.
- 7.
- 8.
9. `value = *newptr; // pointer dereference`

Running Example – w/o IPEA-San

1. `ptr = malloc(size); // object creation`
2. `ptr_base = ptr;`
3. `ptr_bound = ptr + size;`
4. `newptr = ptr + index; // pointer propagation`
5. `newptr_base = ptr_base;`
6. `newptr_bound = ptr_bound;`
- 7.
- 8.
9. `value = *newptr; // pointer dereference`

Running Example – w/o IPEA-San

```
1. ptr = malloc(size); // object creation
2. ptr_base = ptr;
3. ptr_bound = ptr + size;
4. newptr = ptr + index; // pointer propagation
5. newptr_base = ptr_base;
6. newptr_bound = ptr_bound;
7. if ((newptr < newptr_base) || (newptr+sizeof(*newptr) > bound))
8.     abort();
9. value = *newptr; // pointer dereference
```

Running Example – w/o IPEA-San

1. `ptr = malloc(size); // object creation`
 2. `ptr_base = ptr;`
 3. `ptr_bound = ptr + size;`
 4. `newptr = ptr + index; // pointer propagation`
 5. `newptr_base = ptr_base;`
 6. `newptr_bound = ptr_bound;`
 7. `if ((newptr < newptr_base) || (newptr+sizeof(*newptr) > bound))
 abort();`
 9. `value = *newptr; // pointer dereference`
-
- The diagram illustrates the pointer flow and memory access pattern. It shows the initial allocation of memory at address `ptr`, followed by pointer propagation to `newptr`. The range from `newptr` to `value` is highlighted with a red dashed box, indicating a potential access violation if `newptr` is outside its bounds. Two green callout boxes provide context: one for the need to store per-pointer metadata (covering steps 1-3) and another for the need for on-device computation (covering steps 4-9).

Running Example – w/ IPEA-San

1. `ptr = malloc(size); // object creation`
2. `send_to_PC(OP_NEW, ptr, ptr + size, ptr_id);`
- 3.
4. `newptr = ptr + index; // pointer propagation`
5. `send_to_PC(OP_PROP, newptr_id, ptr_id);`
- 6.
7. `send_to_PC(OP_CHK, newptr, sizeof(*newptr), newptr_id);`
- 8.
9. `value = *newptr; // pointer dereference`

Running Example – w/ IPEA-San

1. `ptr = malloc(size); // object creation`
 2. `send_to_PC(OP_NEW, ptr, ptr + size, ptr_id);`
 - 3.
 4. `newptr = ptr + index; // pointer propagation`
 5. `send_to_PC(OP_PROP, newptr_id, ptr_id);`
 - 6.
 7. `send_to_PC(OP_CHK, newptr, sizeof(*newptr), newptr_id);`
 - 8.
 9. `value = *newptr; // pointer dereference`
-
- The diagram illustrates the execution flow and pointer propagation across multiple steps:
- Step 1:** A solid blue arrow points from the `malloc` call to the `ptr` variable.
 - Step 2:** A dashed red arrow points from the `ptr` variable to the `send_to_PC` call.
 - Step 4:** A solid blue arrow points from the `ptr` variable to the `newptr` assignment.
 - Step 5:** A dashed red arrow points from the `newptr` variable to the `send_to_PC` call.
 - Step 7:** A solid blue arrow points from the `newptr` variable to the `send_to_PC` call.
 - Step 9:** A dashed red arrow points from the `newptr` variable to the `*newptr` dereference.
- Annotations in green boxes:
- Tables are maintained on the PC**: Associated with the `send_to_PC` calls in steps 2, 5, and 7.
 - Computation is done on the PC**: Associated with the `OP_PROP` and `OP_CHK` operations in steps 4 and 6.

Metadata Recovery and Maintenance on PC

```
ptr = malloc(size);
send_to_PC(OP_NEW, ptr, size, ptr_id);
newptr = ptr + index;
send_to_PC(OP_PROP, newptr_id, ptr_id);
send_to_PC(OP_CHK, newptr, sizeof(*newptr), \
           newptr_id);
value = *newptr;
```

- Each pointer is associated with a compiler-generated unique ID.
- Instrumentation streams out pointer operations to the PC
- On receiving pointer operations, the IPEA-San plugin on PC recovers the per-pointer metadata.

Metadata Recovery and Maintenance on PC

```
ptr = malloc(size);  
  
newptr = ptr + index;  
  
value = *newptr;
```



Metadata Recovery and Maintenance on PC

```
ptr = malloc(size);
send_to_PC(OP_NEW, ptr, size, ptr_id);
newptr = ptr + index;

value = *newptr;
```

Pointer ID	Tag	Address	Tag
ptr_id	0x100	ptr	0x100
		ptr	0x100
		ptr	0x100

ID-Tag Table Shadow Memory

Metadata Recovery and Maintenance on PC

```
ptr = malloc(size);
send_to_PC(OP_NEW, ptr, size, ptr_id);
newptr = ptr + index;
send_to_PC(OP_PROP, newptr_id, ptr_id);

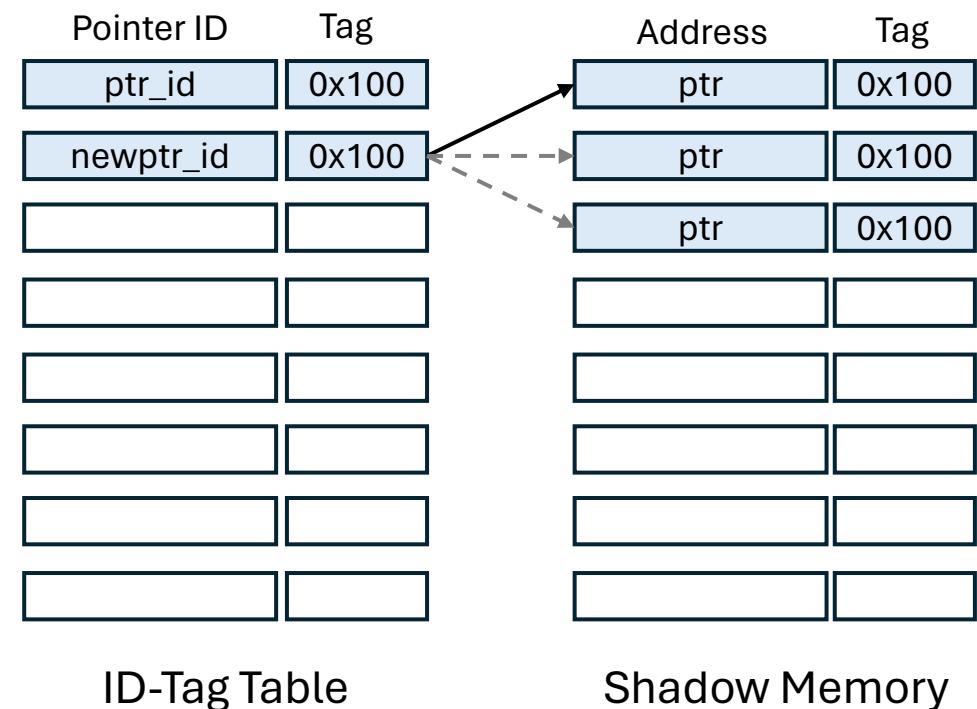
value = *newptr;
```

Pointer ID	Tag	Address	Tag
ptr_id	0x100	ptr	0x100
newptr_id	0x100	ptr	0x100
		ptr	0x100

ID-Tag Table Shadow Memory

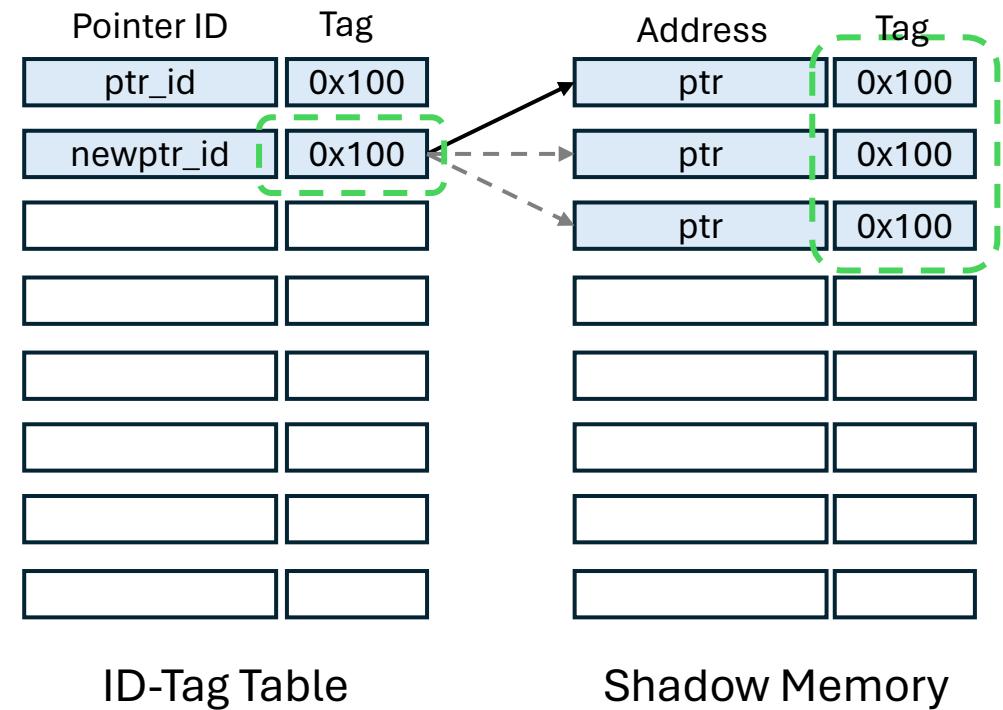
Metadata Recovery and Maintenance on PC

```
ptr = malloc(size);
send_to_PC(OP_NEW, ptr, size, ptr_id);
newptr = ptr + index;
send_to_PC(OP_PROP, newptr_id, ptr_id);
send_to_PC(OP_CHK, newptr, sizeof(*newptr), \
           newptr_id);
value = *newptr;
```



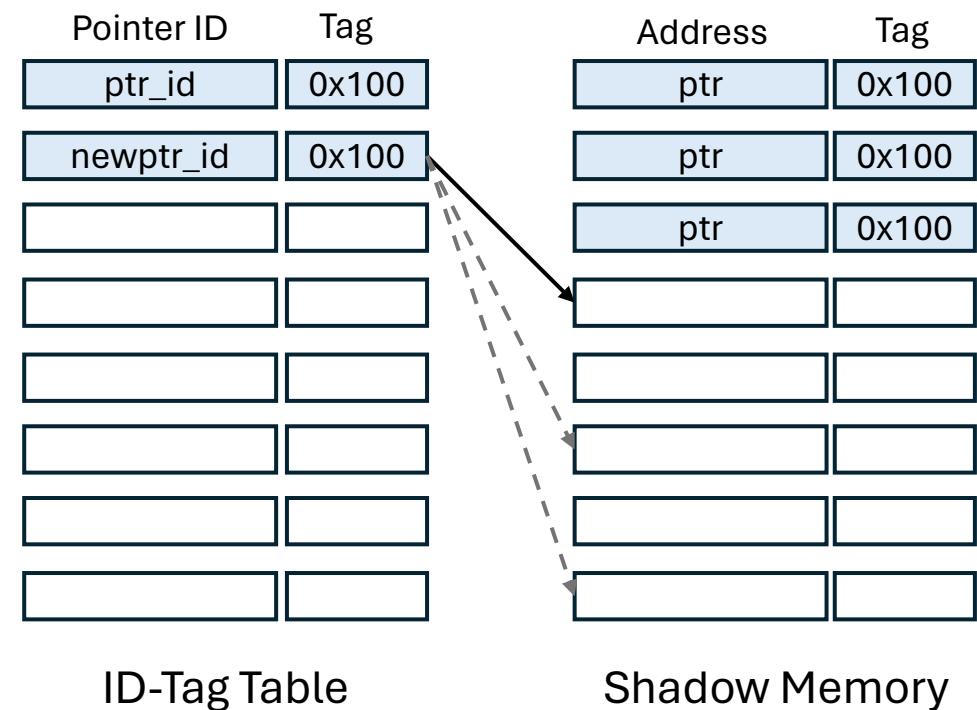
Metadata Recovery and Maintenance on PC

```
ptr = malloc(size);
send_to_PC(OP_NEW, ptr, size, ptr_id);
newptr = ptr + index;
send_to_PC(OP_PROP, newptr_id, ptr_id);
send_to_PC(OP_CHK, newptr, sizeof(*newptr), \
           newptr_id);
value = *newptr;
```



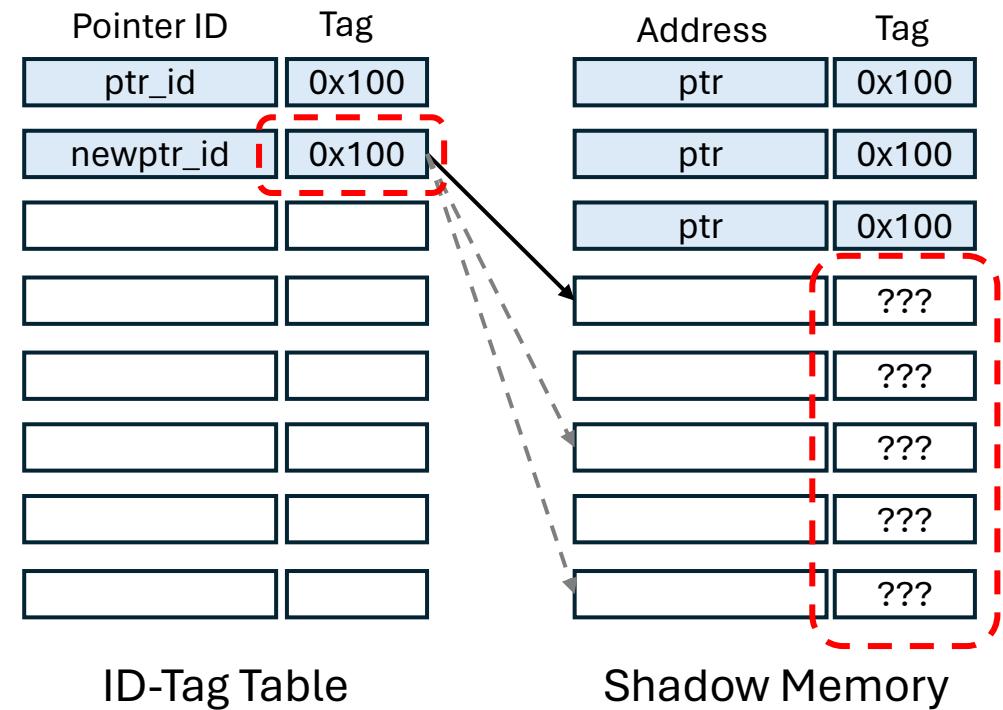
Metadata Recovery and Maintenance on PC

```
ptr = malloc(size);
send_to_PC(OP_NEW, ptr, size, ptr_id);
newptr = ptr + index;
send_to_PC(OP_PROP, newptr_id, ptr_id);
send_to_PC(OP_CHK, newptr, sizeof(*newptr), \
           newptr_id);
value = *newptr;
```



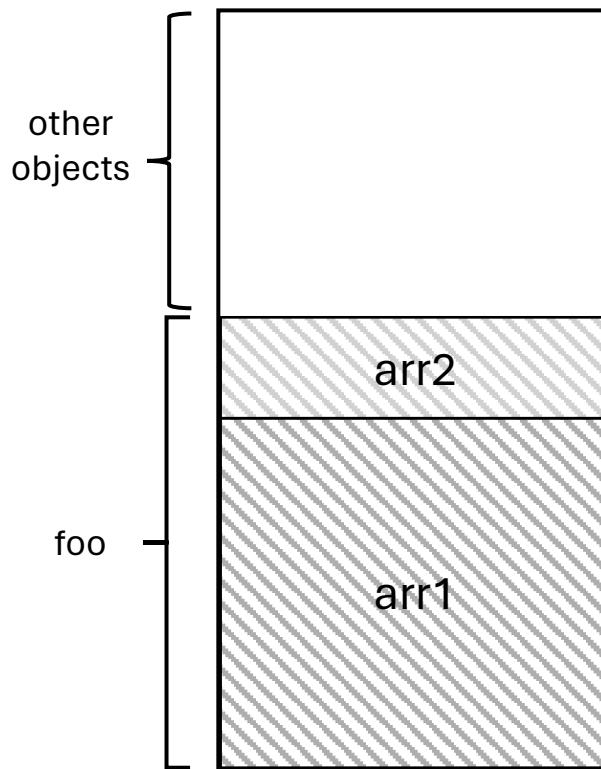
Metadata Recovery and Maintenance on PC

```
ptr = malloc(size);
send_to_PC(OP_NEW, ptr, size, ptr_id);
newptr = ptr + index;
send_to_PC(OP_PROP, newptr_id, ptr_id);
send_to_PC(OP_CHK, newptr, sizeof(*newptr), \
           newptr_id);
value = *newptr;
```



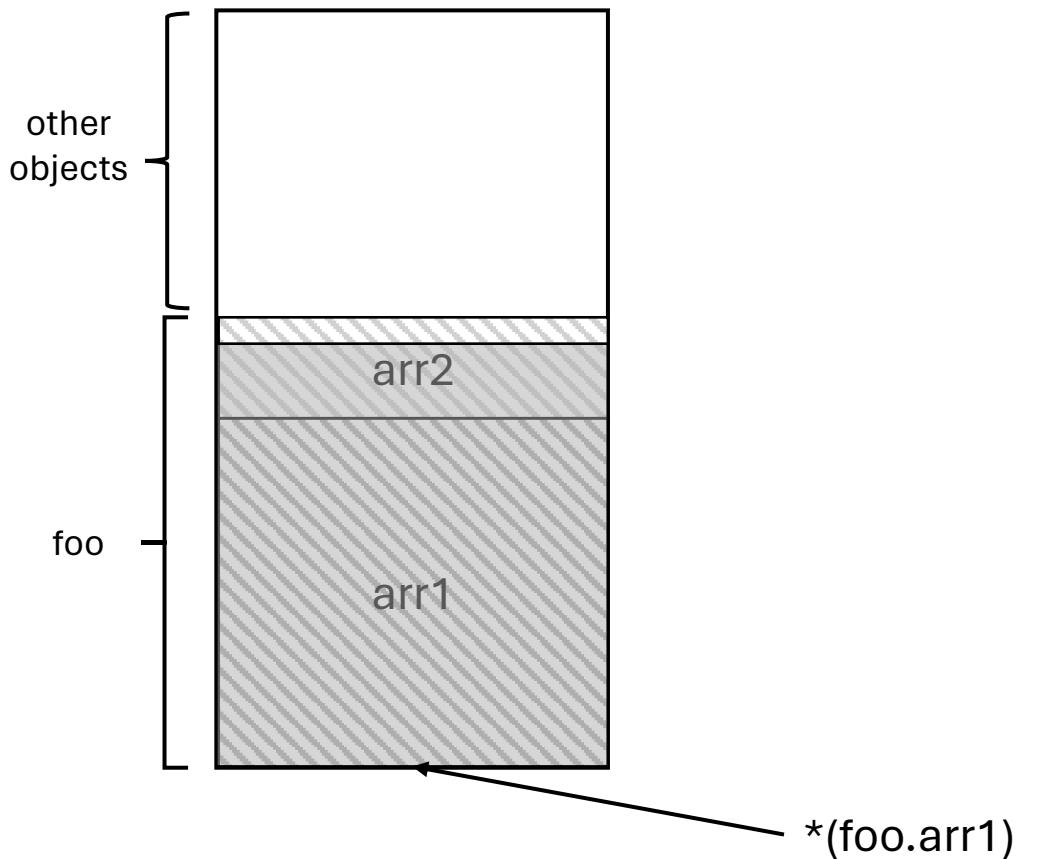
Intra-object Overflow

```
struct foo
{
    char arr1[...];
    char arr2[...];
};
```



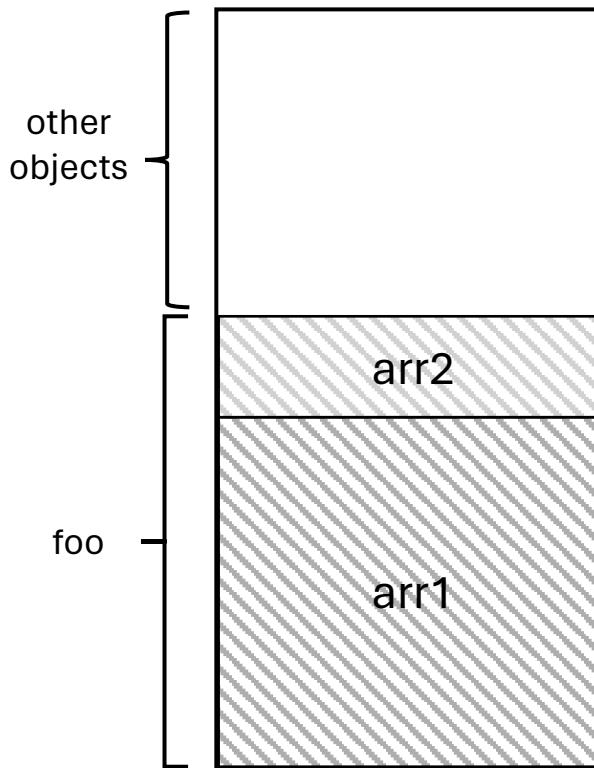
Intra-object Overflow Detection

```
struct foo
{
    char arr1[...];
    char arr2[...];
};
```



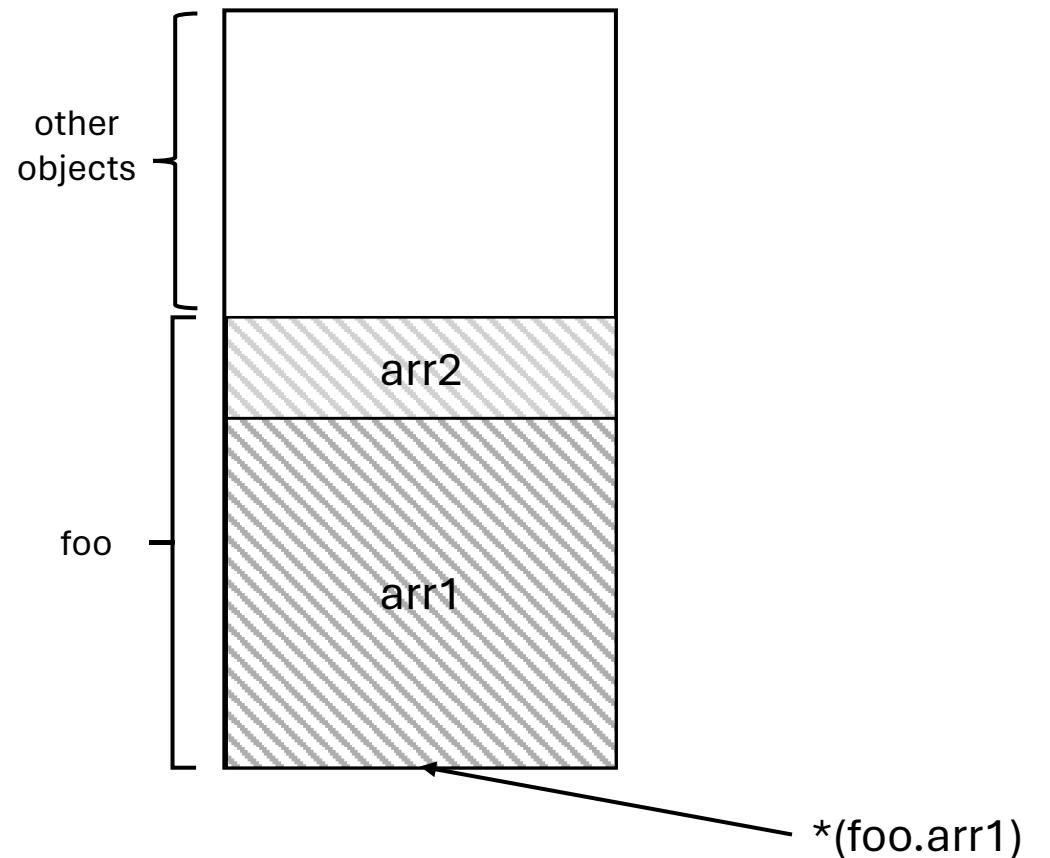
Intra-object Overflow Detection

```
struct foo
{
    char arr1[...];
    char arr2[...];
};
```



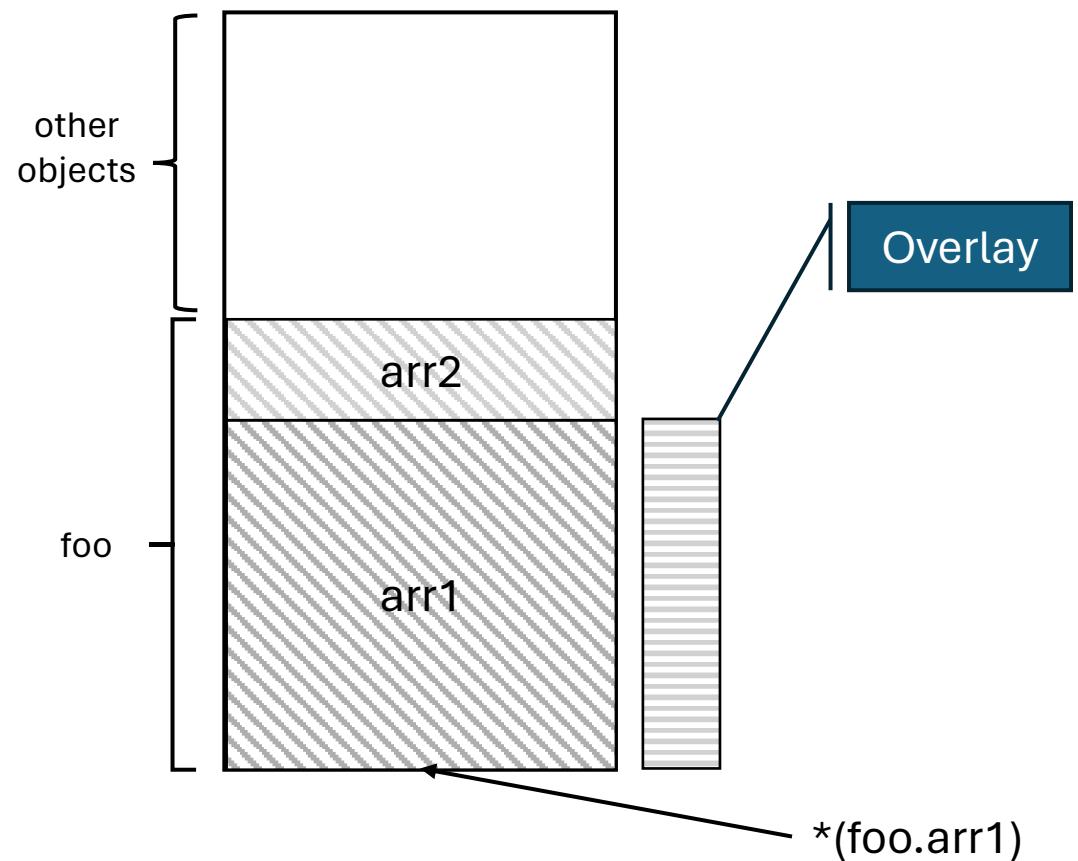
Tag Overlay

```
struct foo
{
    char arr1[...];
    char arr2[...];
};
```



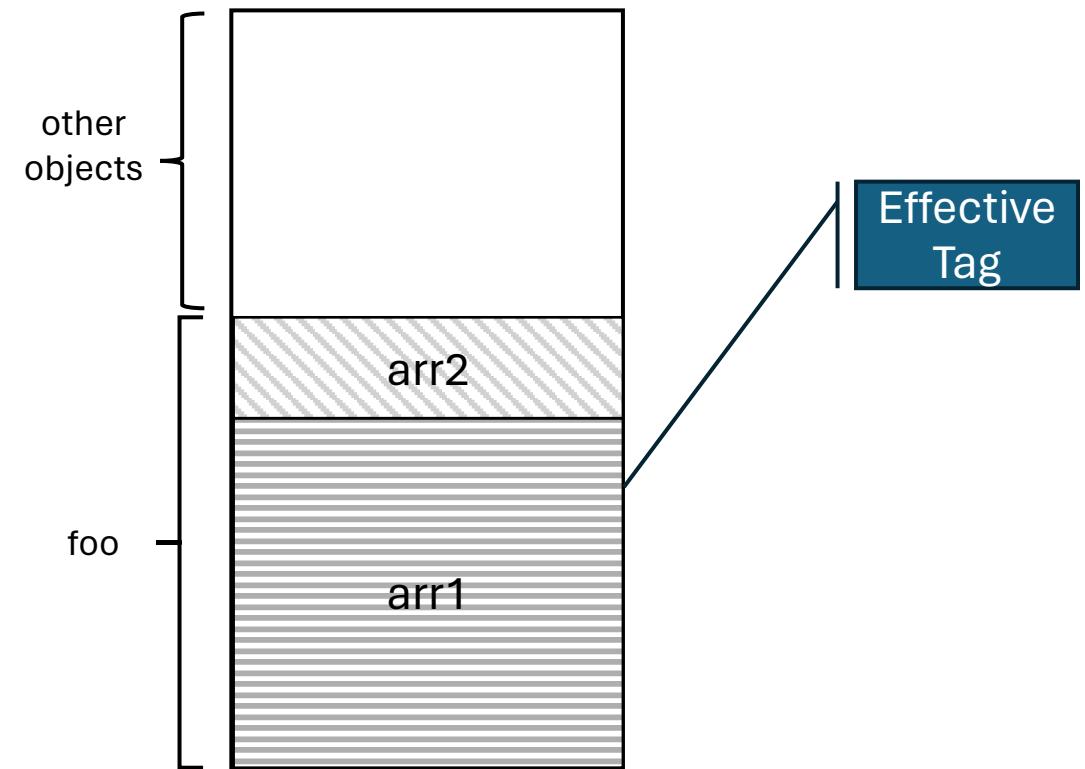
Tag Overlay

```
struct foo
{
    char arr1[...];
    char arr2[...];
};
```



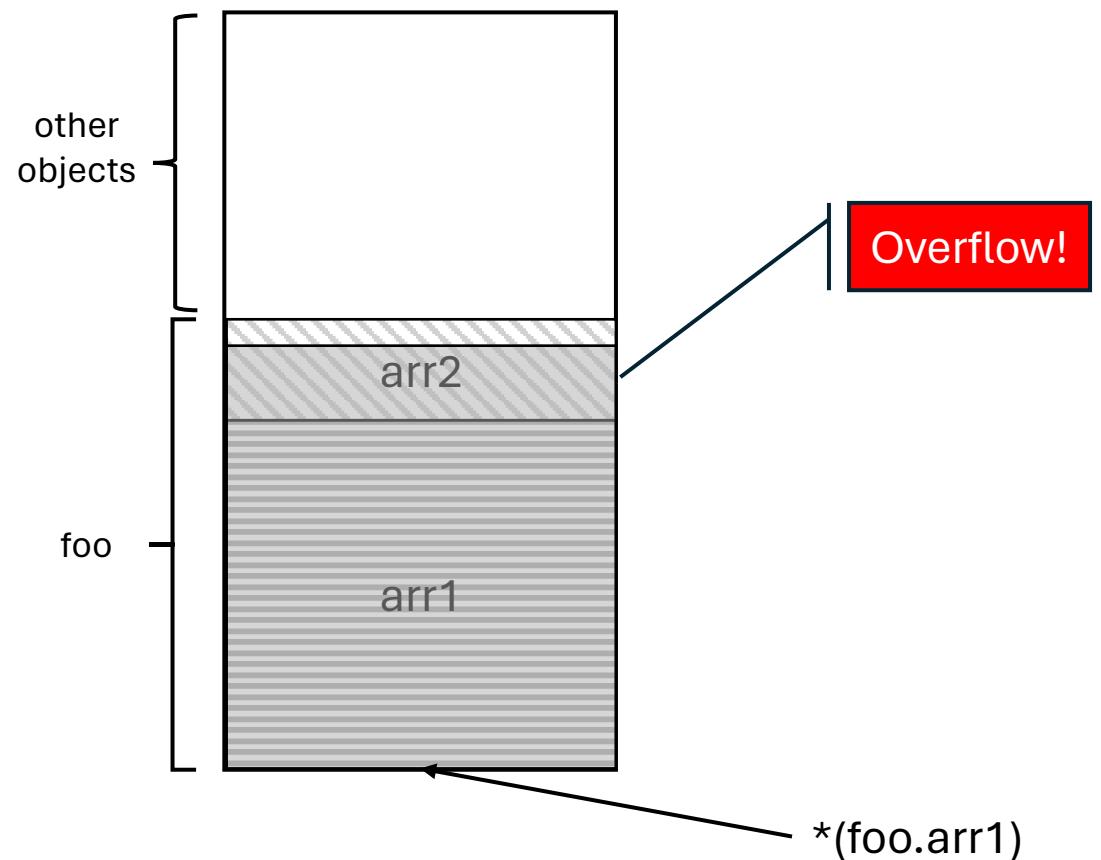
Tag Overlay

```
struct foo
{
    char arr1[...];
    char arr2[...];
};
```



Tag Overlay

```
struct foo
{
    char arr1[...];
    char arr2[...];
};
```



IPEA-Fuzz Plugin

- Firmware Instrumentation
 - A random number is generated at compile-time to identify a basic block.
 - The random numbers are sent to IPEA-Fuzz plugin at run-time.
- IPEA-Fuzz Plugin
 - Maintains a 64KB bitmap shared with AFL to record the number of hits for a particular edge.

Experiment Setup

- Tested prototype boards (7)
 - NXP FRDM-K64F (1MB flash / 256KB SRAM)
 - NXP FRDM-K66F (2MB flash / 256KB SRAM)
 - NXP LPCXpresso55S69 (640KB flash / 320KB SRAM)
 - STM32-NucleoF411R (512KB flash / 128KB)
 - STM32H7B3I-DK (1MB flash / 1MB SRAM)
 - Raspberry Pi Pico (2MB flash / 264KB SRAM)
 - nRF52-DK (192KB flash / 24KB SRAM)
- Debug dongle
 - SEGGER J-Link Edu



Evaluation

- **RQ1:** What kind of memory errors can be captured by IPEA-San?
- **RQ2:** Can IPEA-San reduce resource consumption on MCUs?
- **RQ3:** Can the combination of IPEA-San and IPEA-Fuzz find bugs?

RQ1: IPEA-San Capability

CWE Index	Description	# Tests
CWE 121	Stack-based Buffer Overflow	2,432
CWE 122	Heap-based Buffer Overflow	1,594
CWE 124	Buffer Underwrite	682
CWE 126	Buffer Overread	524
CWE 127	Buffer Underread	682
CWE 415	Double Free	190
CWE 416	Use After Free	118
CWE 476	NULL Pointer Dereference	234
CWE 761	Invalid Heap Pointer Free	152

Selected Testcases in Juliet Testsuite

RQ1: IPEA-San Capability

CWE Index	IPEA-San		ASan	
	False Negative (%)	False Positive (%)	False Negative (%)	False Positive (%)
CWE 121	0	0	1,220 (50.16%)	0
CWE 122	0	0	42 (2.26%)	0
CWE 124	0	0	170 (24.92%)	0
CWE 126	0	0	179 (34.16%)	0
CWE 127	0	0	170 (24.92%)	0
CWE 415	0	0	0	0
CWE 416	0	0	0	0
CWE 476	0	0	0	0
CWE 761	0	0	0	0
Total	0	0	1,781 (26.95%)	0

RQ2: IPEA-San Overhead

Firmware	Baseline					IPEA-San						ASan					
	Flash (Bytes)	Stack (Bytes)	Heap (Bytes)	Global (Bytes)	Total RAM (Bytes)	Flash (x)	Stack (x)	Heap (x)	Global (x)	RTT [†] (Bytes)	Total RAM (x)	Flash (x)	Stack (x)	Heap (x)	Global (x)	Shadow [‡] (Bytes)	Total RAM (x)
aha-compress	4,356	112	-	184	296	1.81	1.14	-	1.00	1,076	4.69	1.94	3.00	-	1.63	2,085	9.19
ctl-stack	4,796	120	816	204	1,140	2.12	1.13	1.00	1.00	1,076	1.96	1.91	2.80	3.13	3.92	2,148	5.12
ctl-string	8,076	104	232	212	548	1.90	1.00	1.00	1.00	1,076	2.97	1.62	2.00	3.03	5.83	2,202	7.94
frac	7,456	240	-	268	508	1.46	1.03	-	1.00	1,076	3.14	1.52	1.60	-	1.64	2,103	5.76
huffbench	4,960	7,824	1,001	188	9,013	2.08	1.01	1.00	1.00	1,076	1.13	2.00	1.13	4.60	2.17	2,099	1.77
sglib-hashtable	6,408	224	800	664	1,688	2.12	1.00	1.00	1.00	1,076	1.64	1.63	2.07	2.14	1.39	2,163	3.12
Geo. Mean	-	-	-	-	-	1.90	-	-	-	-	2.32	1.76	-	-	-	-	4.78
PinLock	20,956	416	-	760	1,176	1.67	1.06	-	1.00	1,076	1.94	3.88	1.81	-	72.33	8,919	54.97
CNC	74,800	344	-	12,708	13,052	1.44	1.17	-	1.00	1,076	1.09	1.37	1.51	-	2.00	3,684	2.27
nRF52-Keyboard	52,428	478	-	6,224	6,702	2.34	1.05	-	1.09	1,076	1.25	1.70	2.17	-	1.80	1,540	2.05
ClockAndWeather	286,184	732	5,456	179,164	185,360	1.21	2.03	1.00	1.00	1,076	1.01	1.30	1.57	1.01	1.84	45,390	2.06
AudioPlayer	101,832	872	-	13,992	14,864	1.56	1.11	-	1.07	1,076	1.14	1.39	2.26	-	4.01	7,272	4.40
WeighScale	24,052	768	-	21,060	21,828	2.79	1.66	-	1.02	1,076	1.01	2.07	2.32	-	2.86	7,754	3.19
HttpServer*	71,932	840	-	65,408	66,248	2.67	1.05	-	1.01	1,076	1.03	-	-	-	-	-	-
U-Disk*	33,356	188	12,168	44,632	56,988	2.75	1.96	1.00	1.00	1,076	1.02	-	-	-	-	-	-
MQTT-Echo*	63,976	632	10,436	59,020	70,088	2.56	1.40	1.00	1.00	1,076	1.02	-	-	-	-	-	-
App-Scheduling	355,108	124	-	7,764	7,888	1.10	1.22	-	1.00	1,076	1.13	1.09	1.42	-	1.32	1,250	1.48
App-Timers	346,816	112	-	7,696	7,808	1.11	1.21	-	1.00	1,076	1.14	1.08	1.42	-	1.32	1,250	1.48
App-IRQs	356,588	108	-	7,800	7,908	1.12	1.17	-	1.00	1,076	1.14	1.07	1.41	-	1.29	1,250	1.45
Geo. Mean	-	-	-	-	-	1.74	-	-	-	-	1.14	1.52	-	-	-	-	3.06

†: RTT buffer plus other metadata.

‡: Shadow memory plus other metadata.

*: Failed to compile using ASan.

RQ3: IPEA-Fuzz Evaluation

Firmware	Target MCU	Time (s)	# Execution	Exec/sec	Paths	Crashes/Hangs
Toy	K64F	7,200	1,990,296	276.43	9	8/0
Expat XML	STM32H7	86,400	1,539,648	17.82	1,098	5/0
JPEGDEC	STM32H7	83,752	805,308	9.61	903	17/4
PNGdec	STM32H7	84,011	848,512	10.10	1,001	22/10
UART	K64F	86,471	537,849	6.22	22	0/0
WiFi	STM32H7	151,212	32,218	0.21	153	1/5
USB Host	K64F	327,910	127,884	0.39	99	22/47
USB Host	STM32H7	347,774	141,277	0.45	96	0/57
emUSB-Host	LPC55S69	122,452	87,619	0.88	106	3/23
microSD	K64F	86,457	30,240	0.35	77	0/6

RQ3: IPEA-Fuzz Evaluation

Firmware	Target MCU	Time (s)	# Execution	Exec/sec	Paths	Crashes/Hangs
Toy	K64F	7,200	1,990,296	276.43	9	8/0
Expat XML	STM32H7	86,400	1,539,648	17.82	1,098	5/0
JPEGDEC	STM32H7	83,752	805,308	9.61	903	17/4
PNGdec	STM32H7	84,011	848,512	10.10	1,001	22/10
UART	K64F	86,471	537,849	6.22	22	0/0
WiFi	STM32H7	151,212	32,218	0.21	153	1/5
USB Host	K64F	327,910	127,884	0.39	99	22/47
USB Host	STM32H7	347,774	141,277	0.45	96	0/57
emUSB-Host	LPC55S69	122,452	87,619	0.88	106	3/23
microSD	K64F	86,457	30,240	0.35	77	0/6

Found 7 zero-day bugs:

- 3 in popular IoT libraries
- 4 in peripheral driver code, including a commercial product

Conclusions

- We present the design and implementation of IPEA firmware analysis framework and two analysis plugins for it.
- They seamlessly integrate into existing firmware development environments, allowing developers to run advanced firmware testing while developing firmware.
- By offloading analysis to the development PCs, the proposed analysis techniques significantly reduce memory overhead compared with solutions that run entirely on MCUs.

Thank you!!

jiameng@uga.edu

<https://github.com/MCUSec/IPEA>