# Private Aggregate Queries to Untrusted Databases

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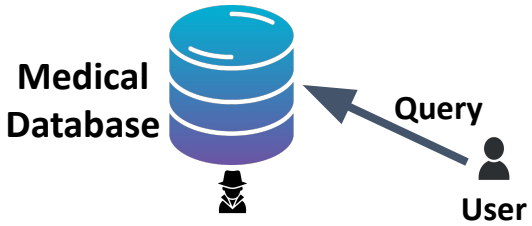
# Aggregate Information Retrieval with Privacy

- Motivation: Data provider can observe all queries run on their database by any user, the computations taking place on the server, and which database rows are scanned




- **Goal:** Retrieve information from an untrusted database without revealing specific queries, even in the presence of *t* colluding database servers

```
SELECT COUNT(user_id)
FROM patients
WHERE is_smoker = 'yes'
AND cancer_flag = 1
```

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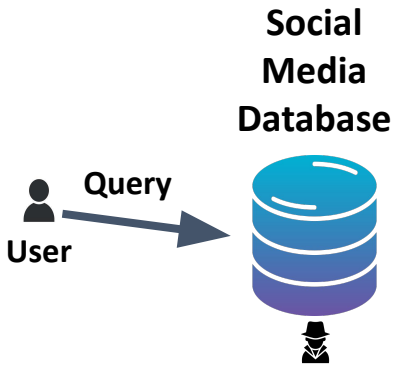


```
SELECT SUM(num_likes)
FROM tweets
WHERE user_id = '20124'
AND date ≤ getdate()
```

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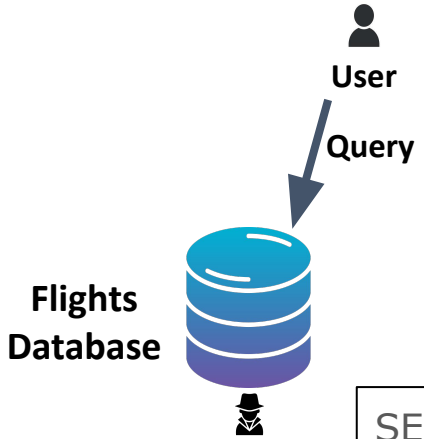
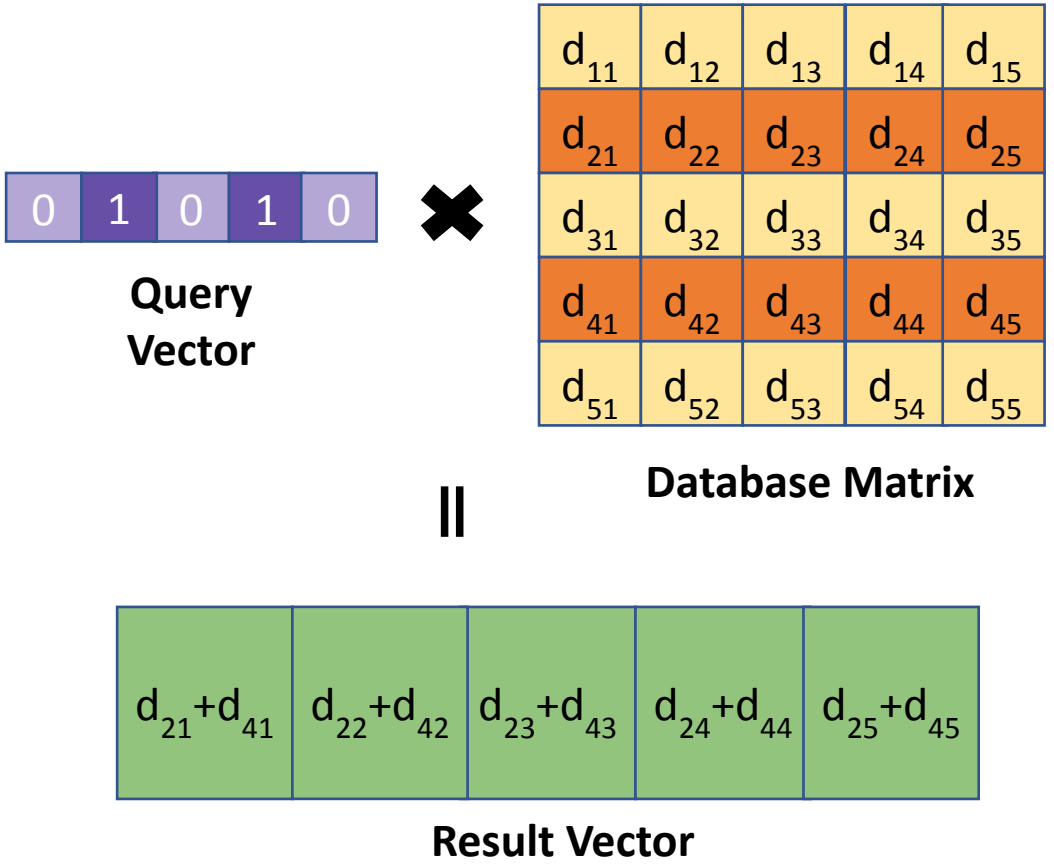



SELECT SUM(***price***)
FROM ***flights***
WHERE ***flight_id*** in
(**'1120', '4268'**)

# Vector Matrix Model

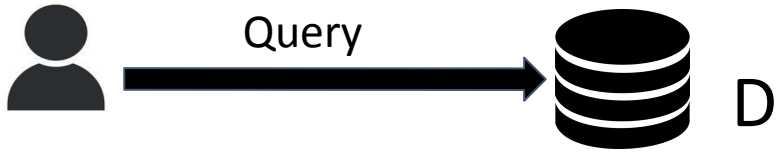
- Database modeled as an *r x s* matrix where *r* corresponds to the number of data blocks (or rows) [*Goldberg, 2007*]
- To fetch the block of data, r-dimensional query vector encoded with a 1 in the i-th position and 0s at every other index
- Product of this query vector with the database matrix produces the desired block of data
- However, this procedure is not private and so we use linear secret sharing



# Making VMM Private for Information Retrieval

- User shares query vector component-wise across servers, share vectors are multiplied with copies of database matrix hosted in each server, and user receives independent products from each server

- User performs component-wise reconstruction using responses received from the servers to obtain desired block of data



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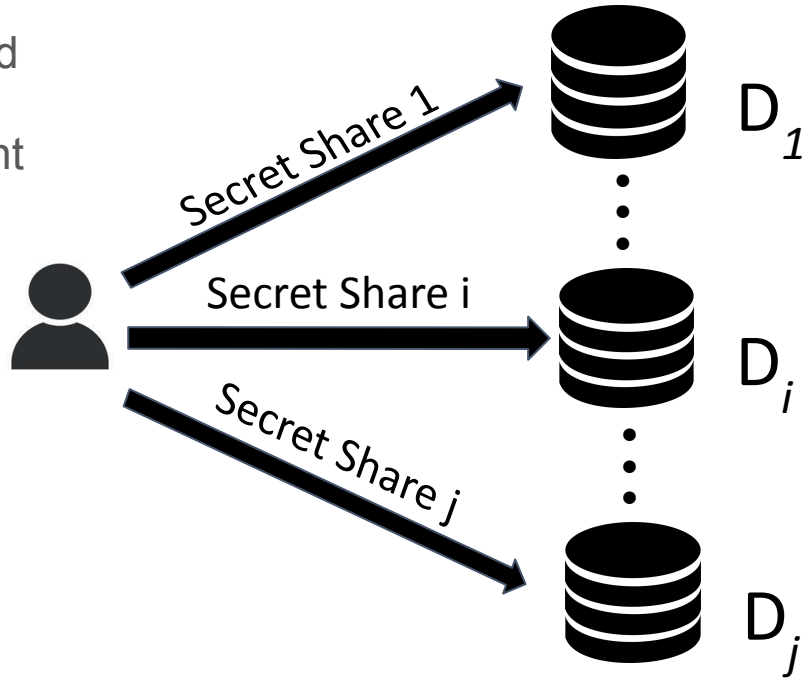

# Proposed PIR: Indexes of Aggregate Queries

- Data structure that maps the database into another matrix, designed to serve specific queries
- Each column corresponds to a row in the database, each row corresponds to a unique value of an attribute in the database
- Multiple indexes of queries can be batched together if dimensions same [*Hafiz-Henry, 2017*]

D =

| Hospitalization_ID | Patient_ID | Admit_Date | Gender_ID | Days_Hospitalized | State_ID |
|---|---|---|---|---|---|
| 1 | 1 | 01-02-2022 | 1 (Male) | 10 | 2 (OR) |
| 2 | 2 | 01-04-2022 | 1 (Male) | 2 | 1 (CA) |
| 3 | 3 | 08-06-2022 | 2 (Female) | 14 | 3 (WA) |
| 4 | 1 | 07-23-2022 | 1 (Male) | 2 | 2 (OR) |
| 5 | 3 | 09-01-2022 | 2 (Female) | 7 | 3 (WA) |
| 6 | 4 | 05-14-2022 | 3 (Other) | 2 | 1 (CA) |

$\prod_{\text{patient}} :=$

| | | | | | | |
|---|---|---|---|---|---|---|
| patient 1 | 1 | 0 | 0 | 1 | 0 | 0 |
| patient 2 | 0 | 1 | 0 | 0 | 0 | 0 |
| patient 3 | 0 | 0 | 1 | 0 | 1 | 0 |
| patient 4 | 0 | 0 | 0 | 0 | 0 | 1 |

# Sample Query

D =

| Hospitalization_ID | Patient_ID | Admit_Date | Gender_ID | Days_Hospitalized | State_ID |
|---|---|---|---|---|---|
| 1 | 1 | 01-02-2022 | 1 (Male) | 10 | 2 (OR) |
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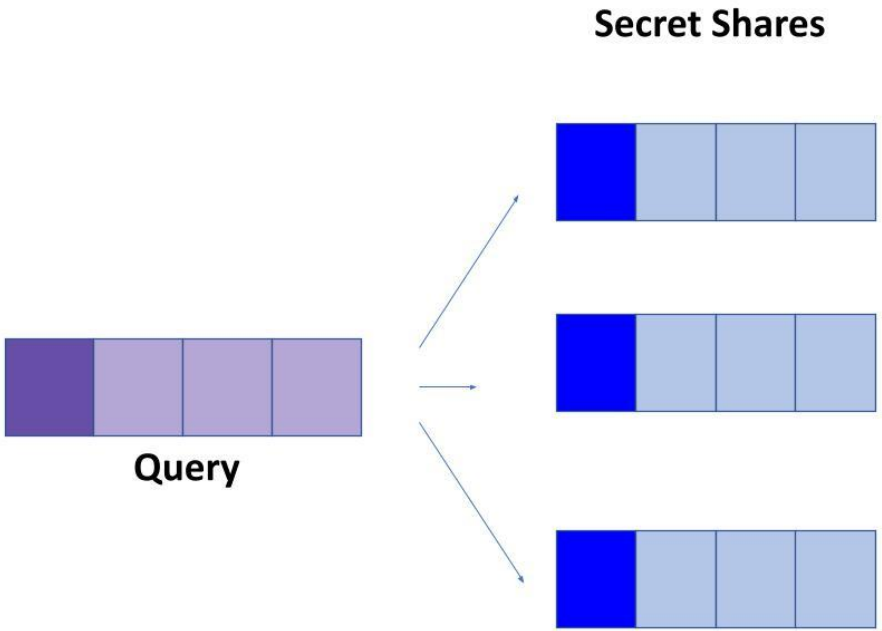
| patient 4 | 0 | 0 | 0 | 0 | 0 | 1 |

```
SELECT SUM(Days_Hospitalized)
From D
WHERE Patient_ID = 1
```

| | | | |
|---|---|---|---|
| 1 | 0 | 0 | 0 |

Query Vector

# Protocol Schematic



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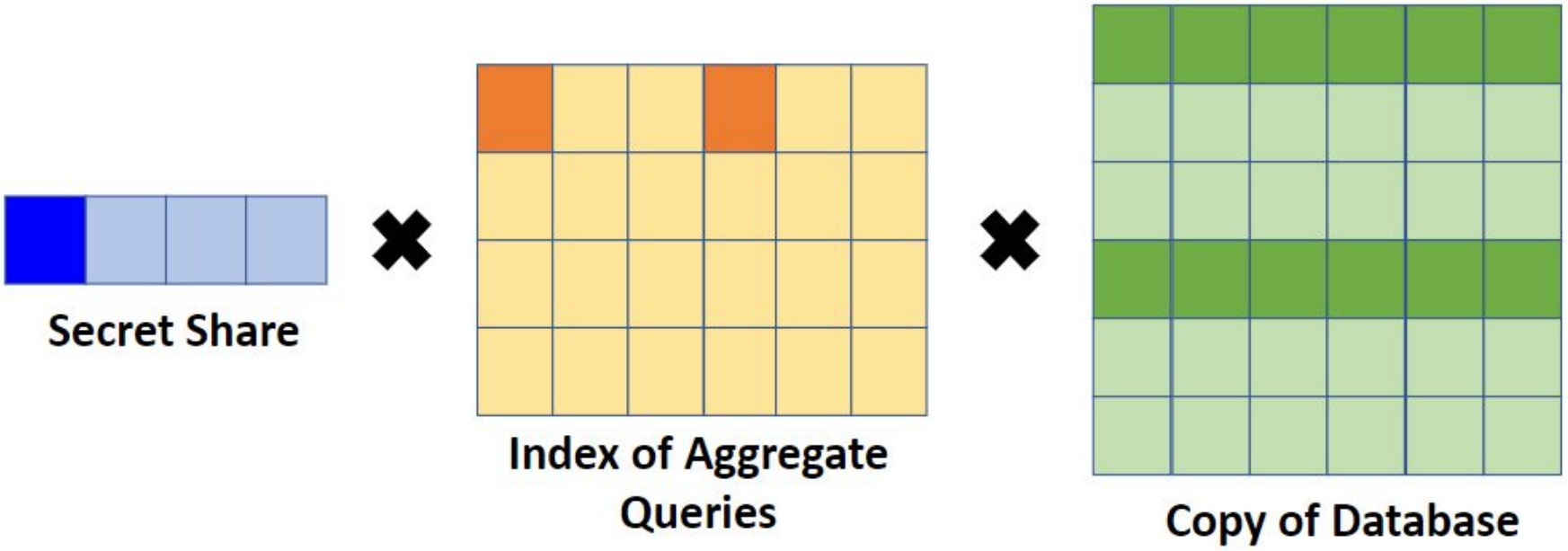

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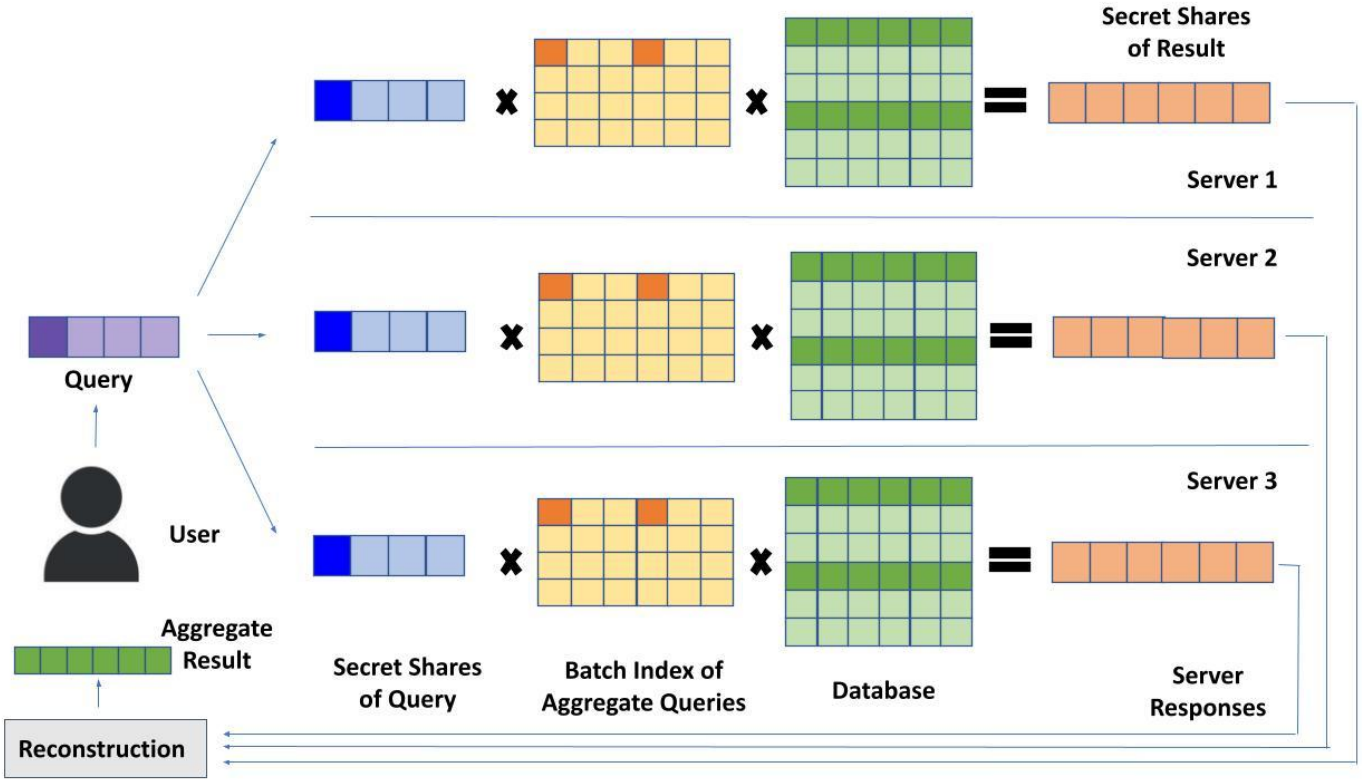

# Case Studies

X (Formerly Twitter)

- Scraped 1,004,129 tweets with politically relevant hashtags such as ‘USElections’, ‘Trump’, ‘Biden’
- 2 indexes of queries batched to serve queries about like counts and retweet counts, each index of query of dimension 333,286 × 1,004,129
- Each row in the index of queries corresponds to a unique user in the scraped database

MIMIC 3

- Clinical dataset of hospitalization records
- First set batches 4 indexes of queries to serve 4 different queries, each index of query matrix is of dimension 4 x 58,976, with each row corresponding to a different value of admission type
- Second set batches 2 indexes of queries to serve 2 queries, each index of query is of dimension 1,400 × 4,156,450, with each row corresponding to a different patient

# Case Studies

X (Formerly Twitter)

SELECT SUM(**like_count**) FROM **twitter_data** WHERE **user_id** = **‘100012’**

SELECT COUNT(*) FROM **twitter_data** WHERE **user_id** = **‘100012’** AND **no_retweets** = **0**

MIMIC 3

SELECT SUM(**hospitalization_duration**) FROM **admissions** WHERE **subject_id** = **‘100012’** AND **admission_type** = **‘EMERGENCY’**

SELECT COUNT(*) FROM **admissions** WHERE **admission_type** = **‘URGENT’**

# Case Study Results

| Case Study Database | Index of Aggregate Queries for | Index Matrix Dimension | Index Generation Time (secs) | Batching Time for Multiple Indexes (mins) | Additional Data Structure Storage Size (MiB) | VSpM Throughput on GPU (clients/sec) | Server Response Generation Time (secs) | | |
|---|---|---|---|---|---|---|---|---|---|
| | | | | | | | All Attributes | Essential Attributes | Baseline |
| **MIMIC 3** | Admission Type | 4 x 58,976 | 0.06 | 0.002 | 0.76 | 20,534.12 | **0.11** | **0.04** | **0.11** |
| | Ethnicity | | 0.39 | | | | | | |
| | Latest Admission | | 0.98 | | | | | | |
| | Oldest Admission | | 0.95 | | | | | | |
| | Dosage | 1,400 x 4,156,450 | 2.12 | 0.101 | 34.34 | 4,412.80 | **0.26** | **0.10** | **7.37** |
| | Stay Duration | | 2.03 | | | | | | |

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# Server Response Generation on Larger Databases

| DB Size | Case Study | Records | Record Size | Protocol | $\mathbf{GF}(2^8)$ | $\mathbf{GF}(2^{16})$ | $-mZZ_p - w128$ | $-mZZ_p - w256$ |
|---|---|---|---|---|---|---|---|---|
| 40 GiB | Twitter Filt. | 1,004,129 | 41.8 KiB | Baseline | 16.8 | 32.1 | 1006.1 | 580.6 |
| | | | | This work | 0.1 | 0.2 | 7.7 | 4.9 |
| | MIMIC 3 Filt. | 4,156,450 | 10.5 KiB | Baseline | 17.5 | 33.6 | 1010.6 | 589.3 |
| | | | | This work | 0.6 | 1.1 | 36.6 | 21.2 |
| 64 GiB | Twitter Filt. | 1,004,129 | 67.0 KiB | Baseline | 27.5 | 51.0 | 1754.1 | 988.8 |
| | | | | This work | 0.2 | 0.3 | 12.2 | 7.6 |
| | MIMIC 3 Filt. | 4,156,450 | 16.5 KiB | Baseline | 27.8 | 51.7 | 1703.8 | 981.3 |
| | | | | This work | 0.9 | 1.5 | 49.9 | 33.2 |

Response times for all modulus bit sizes are in seconds

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# Takeaways

- Novel framework that augments conventional IT-PIR protocols (e.g., Goldberg's IT-PIR) with aggregate queries
  - Constructions of effective indexes of aggregate queries comprising new standard aggregate vector

- Simulated real-world applications to benchmark performance and scalability of proposed PIR scheme with aggregate queries

- Efficient implementation of our framework on GPU can achieve fast query response time while assuring the privacy of aggregate queries

# Thank You