

# Faster and Better: Detecting Vulnerabilities in Linux-based IoT Firmware with Optimized Reaching Definition Analysis

Zicong Gao<sup>+</sup>, Chao Zhang<sup>\*</sup>, Hangtian Liu, Wenhou Sun, Zhizhuo Tang, Lihui Jiang, Jianjun Chen, and Yong Xie



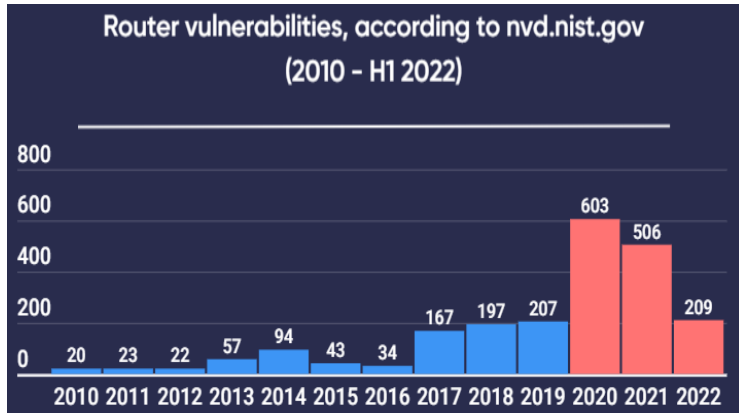
# IoT Devices

- The number of IoT devices has reached 15.14 billion by the end of 2023
- Rich application scenarios from life to production
  - ◆ e.g. Smart Devices, Wearables, Webcams, Connected Vehicles, Industrial control system



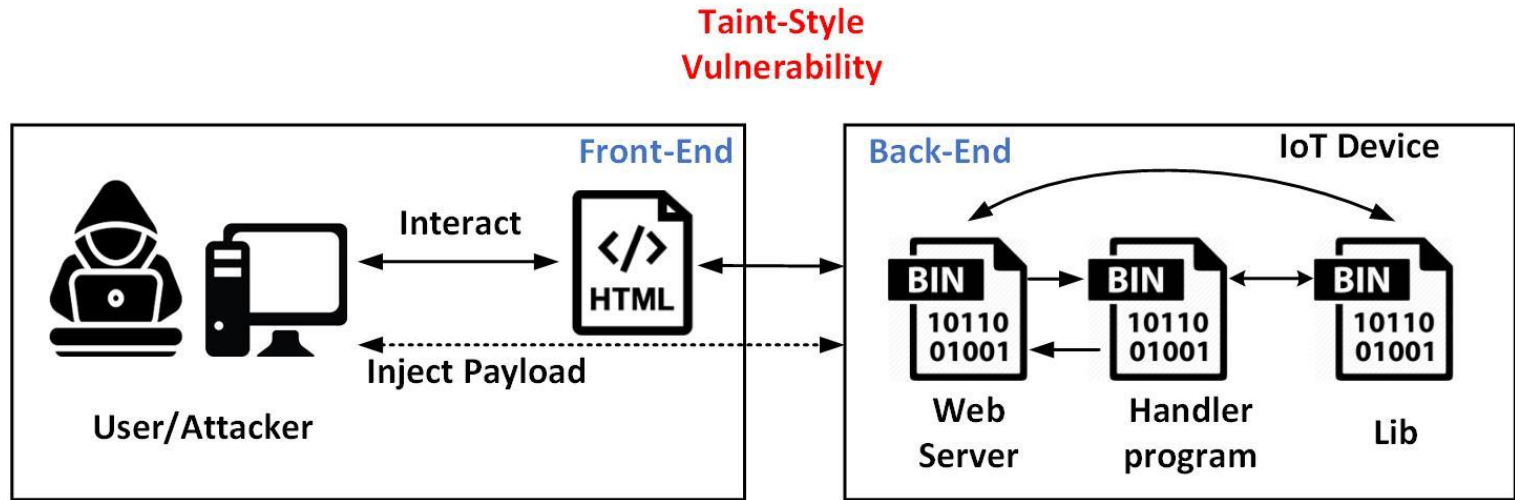
# IoT Device Vulnerabilities

- IoT devices suffer from the serious cyber threats
  - ◆ Network devices (e.g. router, webcam, firewall) are the most commonly attacked IoT devices
  - ◆ Vulnerabilities in network security equipment have extremely serious impacts



# Taint-style Vulnerabilities in IoT

- Vulnerabilities, especially **taint-style vulnerabilities**, are significant security threats to IoT devices



Threat model of taint-style vulnerabilities in IoT scenarios

How to detect taint-style vulnerabilities in  
IoT devices?

# Proposed Method

- Dynamic solutions

- Fuzzing

Pros: accurate, high true positive rate

Cons: requiring emulation, difficult to explore deep paths

- Static solutions

- Taint analysis

Pros: scalability, high coverage

Cons: high false positive rate, heavyweight symbolic execution

# Motivation Example

- Limitation of existing works

- SaTC & KARONTE fail to alert
- Reason: functions containing source & sink points are only indirectly called functions or library functions, neglected by common CFG construction strategies
- Example: the function `ej_hwdpi_monitor_info` is not in the CFG and the `libbwdpi_sql.so` is not analyzed either

main()->handle\_request() looks up table mine\_handler

appGet.cgi()->do\_ej() looks up table ej\_handler

```
//httpd
main()->handle_request()->mime_handler
mime_handler
0x9C2C4 DCD aAppGetCgi ;"appGet.cgi*"
0x9C2C8 DCD aTextHtml ;"text/html"
0x9C2CC DCD aCacheControl ;"Cache-Control:..."
0x9C2D0 DCD sub_283D4 ;"handler address"
0x9C2D4 DCD appGet.cgi ;"handler name"
```

Subgraph-1

```
// httpd
appGet.cgi()->do_ej()->ej_handler
ej_handler
0x9B924 DCD BwdpIM ;"ej_bwdpi_monitor_info"
0x9B928 DCD sub_2DE14 ;"function address"
```

Subgraph-2

```
//httpd
1 int ej_bwdpi_monitor_info(webs wp, int argv){
2 int retval[5];
3 retval[1] = argv;
4 retval[0] = 0;
5 char *type = websGetVar("type", wp, "local");
6 if(!type)
7 type = "";
8 char *event = websGetVar("event", wp, "login");
9 if(!event)
10 event = "";
11 bwdpi_monitor_info(type, event, retval, wp);
12 return retval[0];
13 }
```

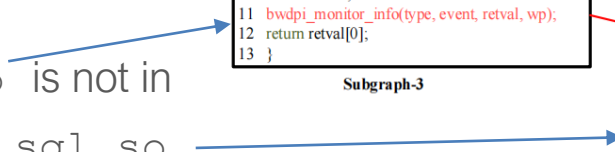
Subgraph-3

```
//httpd
1 char *websGetVar(char *var, webs wp, char *defaultGetValue){
2 char *result;
3 sym *sp = symLookup(wp->cgiVars, var);
4 if(!sp) return defaultGetValue;
5 result = sp->content.value.string;
6 return result;
7 }
```

Subgraph-4

```
//libbwdpi_sql.so
1 int bwdpi_monitor_info(char *type, char *event, int retval, webs wp){
2 char v74[1056];
3 if(f_exists("log")>0){
4 sprintf(v74, 0x400,
5 "echo event=%s >> log", event);
6 system(v74);
7 .....
8 }
```

Subgraph-5



# Motivation Example

- Limitation of existing works

- Symbolic execution-based taint tracking is ...

- **time consuming**: SaTC costs 0.5h~30h per sample, and the analysis time will increase by 2 to 3 times when libraries are included.....

- **not practicable in real-world**: It faces problems such as state explosion, path explosion, and constraint solving complexity.....



# Motivation Example

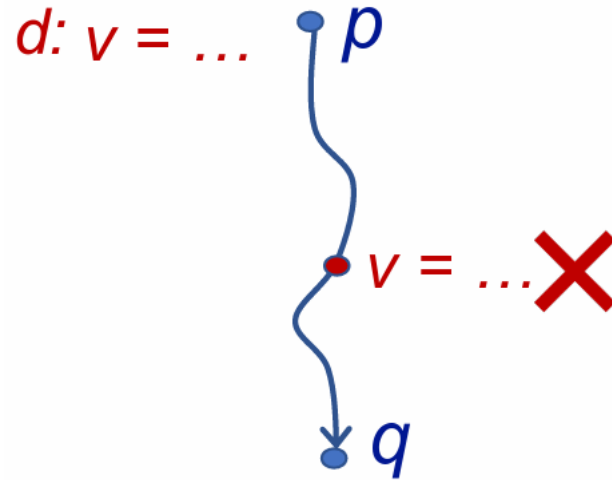
- Limitation of existing works
  - Symbolic execution-based taint tracking is ...
    - **time consuming**: SaTC costs 0.5h~30h per sample, and the analysis time will increase by 2 to 3 times when libraries are included.....
    - **not practicable in real-world**: It faces problems such as state explosion, path explosion, and constraint solving complexity.....

**More often than not, static analysis should be fast and productive**

# Our solution

- Reaching definition analysis

A definition  $d$  of a variable  $v$  at program point  $p$  reaches a point  $q$  if there is a path from  $p$  to  $q$  such that  $d$  is not “killed” along that path



# Our solution

- Reaching definition analysis

A definition  $d$  of a variable  $v$  at program point  $p$  reaches a point  $q$  if there is a path from  $p$  to  $q$  such that  $d$  is not “killed” along that path

- RDA-based taint analysis

A definition  $d$  of a **taint variable**  $v$  at **source**  $p$  reaches **sink**  $q$  if there is a path from  $p$  to  $q$  such that  $d$  is not “killed” along that path

# Our solution

- RDA-based taint analysis
  - Definitions to taint variable `type` and `event` at sources line 5 & 8 assigned by `websGetVar` reaches sink function `system()`
  - Definition value of variable `v74` is related to `type/event` and violates the vulnerability rule

```
//httpd
1 int ej_bwdpi_monitor_info(webs wp, int argv){
2 int retval[5];
3 retval[1] = argv;
4 retval[0] = 0;
5 char *type = websGetVar("type", wp, "local");
6 if(!type)
7 type = "";
8 char *event = websGetVar("event", wp, "login");
9 if(!event)
10 event = "";
11 bwdpi_monitor_info(type, event, retval, wp);
12 return retval[0];
13 }
```

Subgraph-3

```
//httpd
1 char *websGetVar(char *var, webs wp, char *defaultGetValue){
2 char *result;
3 sym *sp = symLookup(wp->cgiVars, var);
4 if(!sp) return defaultGetValue;
5 result = sp->content.value.string;
6 return result;
7 }
```

Subgraph-4

```
//libbwdpi_sql.so
1 int bwdpi_monitor_info(char *type, char *event, int retval, webs wp){
2 char v74[1056];
3 if(f_exists("log")>0){
4 snprintf(v74, 0x400,
5 "echo event=%s >> log", event);
6 system(v74);
7 .....
8 }
```

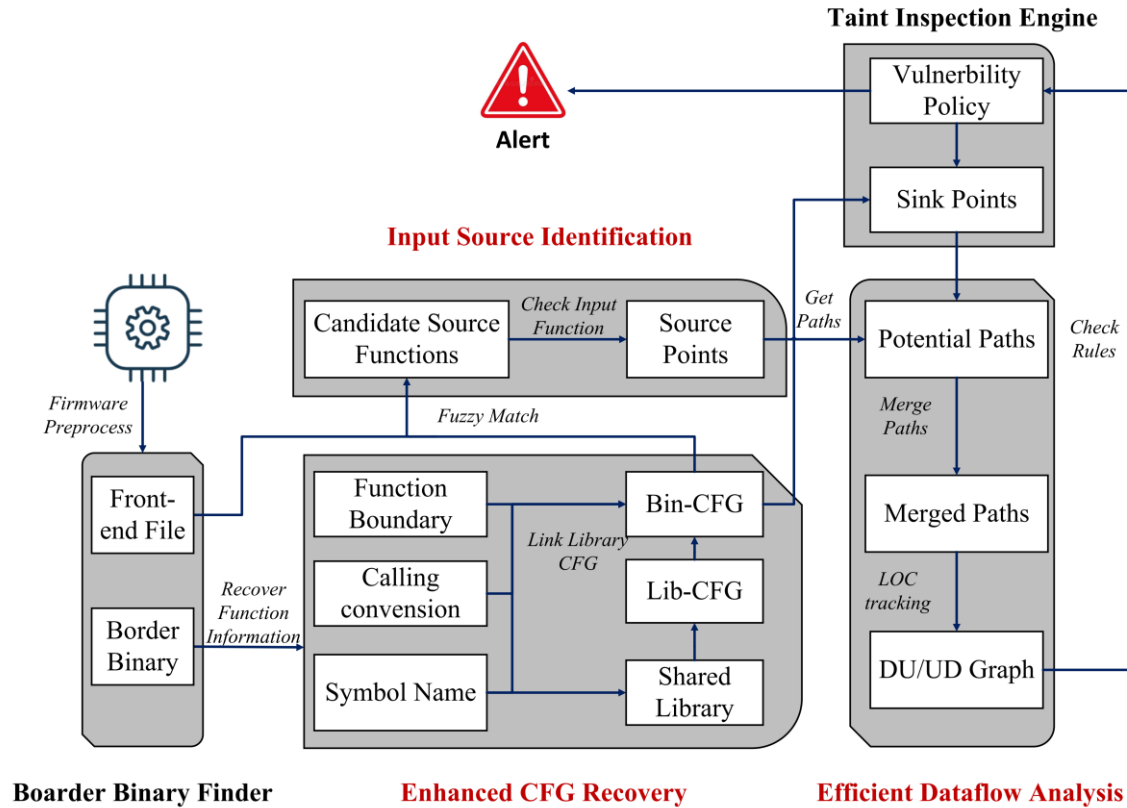
Subgraph-5

# Challenges

- **1. Comprehensive CFG Recovery**
  - Functions only invoked by indirect calls are difficult to identify
  - Connecting binary CFG and library CFG increases analysis efforts
- **2. Precise Source Point Identification**
  - Manually specifying source functions requires expert knowledge and customization
  - Pattern-based string matching methods cannot utilize semantic information, missing some potential source points
- **3. Efficient Taint Tracking**
  - Massive paths between source/sink points can lead to path explosion in RDA

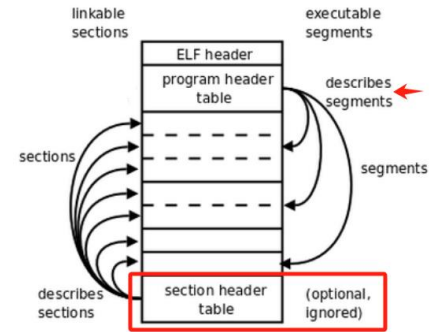
HermeScan

# Architecture



# Enhanced CFG Recovery

- Function boundary
  - Full binary linear scan by dividing function boundaries according to function prologue
- Symbol table
  - When the section header table is stripped, locate the address of symbol table from metadata in `PT_DYNAMIC` segment
- Calling conventions
  - Aggressive but complete recovery strategy by setting default CC for each function
- Connect the Bin-CFG with the Lib-CFG





# Source Input Identification

- Fuzzy Matching

- Consider the **word form similarity** and **semantic similarity** of keywords appearing in the front and back ends. Regard functions reference these matched strings as candidate sources

- The normalized edit distance is used to calculate the FormatSim

$$\text{FormatSim}(S1, S2) = 1 - \text{Edit}(S1, S2) / (L(S1) + L(S2))$$

e.g. hostname\_1.1  hostname\_%s

- The BERT model is used to calculate the semantic similarity of two strings

$$\text{SemanticSim}(S1, S2) = \begin{cases} \text{Cosine}(S1, S2), & \text{Others} \\ 0, & \frac{\text{LCS}(S1, S2)}{\text{Min}(L(S1), L(S2))} < \theta \end{cases}$$

e.g. “request from %s is banned for security”  “sec\_ip\_ban”

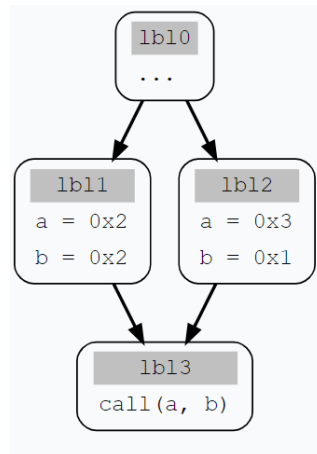
# Source Input Identification

- Candidate Function Checking
  - Goal: Remove infeasible candidate functions and mark parts of the function arguments or return values as taint sources
  - Idea: Check whether the parameters would receive values from external input and whether the return values would be used by following operations

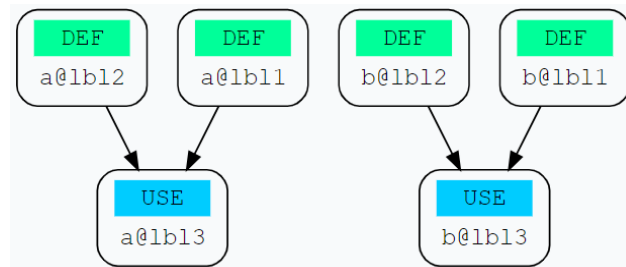
```
//httpd
1 int ej_bwdpi_monitor_info(webs wp, int argv){
2   int retval[5];
3   retval[1] = argv;
4   retval[0] = 0;
5   char *type = websGetVar("type", wp, "local");
6   if(!type)
7     type = "";
8   char *event = websGetVar("event", wp, "login");
9   if(!event)
10    event = "";
11   bwdpi_monitor_info(type, event, retval, wp);
12   return retval[0];
13 }
```

# Efficient Dataflow Analysis

- Lightweight, context-sensitive, on-demand, interprocedural analysis
  - Lightweight: RDA-based instead of symbolic execution-based taint tracking
  - Context-Sensitive: considers context information to enable fine-grained dataflow analysis
  - On-demand:
    - 1) Only step into functions with tainted parameters for interprocedural analysis
    - 2) Use summary for common library functions



Control flow graph



Def-use graph

# Efficient Dataflow Analysis

- Lightweight, context-sensitive, on-demand, inter-procedural analysis

---

```
1 // ssi
2 void setup_wizard_mydlink(int a1){
3     char *v4;
4     v4 = getenv("sys_service");
5     updown_services(0, v4);
6     post2nvram(a1);
7     response_page = get_response_page();
8 }
9 int updown_services(int mode, char *sys_service){
10     if(mode) return func2(sys_service);
11     return func3(sys_service);
12 }
13 int func2(char *a1){
14     char buf[1028];
15     if (a1 && *a1){
16         strcpy(buf, a1);
17     }
18 }
```

---

Mark variable v4 as a taint source

# Efficient Dataflow Analysis

- Lightweight, context-sensitive, on-demand, inter-procedural analysis

---

```
1 // ssi
2 void setup_wizard_mydlink(int a1){
3     char *v4;
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9 int updown_services(int mode, char *sys_service){
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11     return func3(sys_service);
12 }
13 int func2(char *a1){
14     char buf[1028];
15     if (a1 && *a1){
16         strcpy(buf, a1);
17     }
18 }
```

---

Step into the `updown_service`  
function

# Efficient Dataflow Analysis

- Lightweight, context-sensitive, on-demand, inter-procedural analysis

---

```
1 // ssi
2 void setup_wizard_mydlink(int a1){
3     char *v4;
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11     return func3(sys_service);
12 }
13 int func2(char *a1){
14     char buf[1028];
15     if (a1 && *a1){
16         strcpy(buf, a1);
17     }
18 }
```

---

Skip these two functions

# Efficient Dataflow Analysis

- Lightweight, context-sensitive, on-demand, inter-procedural analysis

---

```
1 // ssi
2 void setup_wizard_mydlink(int a1){
3     char *v4;
4     v4 = getenv("sys_service");
5     updown_services(0, v4);
6     post2nvram(a1);
7     response_page = get_response_page();
8 }
9 int updown_services(int mode, char *sys_service){
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11     return func3(sys_service);
12 }
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15     if (a1 && *a1){
16         strcpy(buf, a1);
17     }
18 }
```

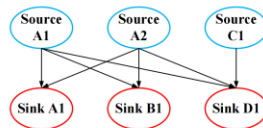
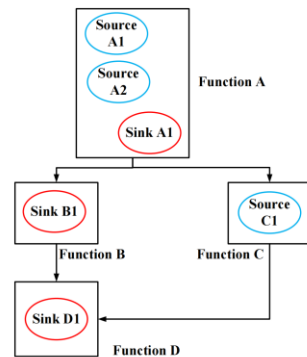
Apply function summary

---

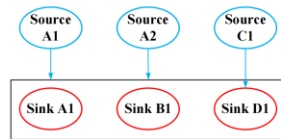
# Efficient Dataflow Analysis

- Path merging strategy

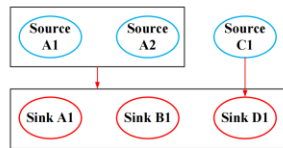
- ◆ Leverages the path-insensitive feature of RDA
- ◆ Multi-source taint:  
taint each source point with a different label  
when a function contains multiple source points
- ◆ Multi-sink observation:  
all sinks in a reachability call graph can be  
observed in one pass of RDA analysis



- Before Merging:
- 1) Source A1->Sink A1
  - 2) Source A2->Sink A1
  - 3) Source A1->Sink B1
  - 4) Source A2->Sink B1
  - 5) Source A1->Sink D1
  - 6) Source A2->Sink D1
  - 7) Source C1->Sink D1



- SaTC Strategy:
- 1) Source A1->Sink A1, Sink B1, Sink D1
  - 2) Source A2->Sink A1, Sink B1, Sink D1
  - 3) Source C1->Sink D1



- HermeScan Strategy:
- 1) Source A1, A2->Sink A1, B1, D1
  - 2) Source C1->Sink D1



# Evaluation

- Q1: How well does HermeScan **find vulnerabilities on real-world devices**? How effective is it **compared to state-of-the-art tools**?
- Q2: How does the optimization of **control flow recovery** contribute to the vulnerability detection of HermeScan?
- Q3: Can HermeScan's **input source identification** make the analysis more accurate? How does it work?
- Q4: Can HermeScan's **path merging strategy** alleviate the path explosion problem?

# Dataset

- 0-day dataset
  - ◆ 30 samples
  - ◆ 8 vendors and 19 series
  - ◆ architecture: ARM32, ARM64, MIPSEL, and MIPSEB
- N-day dataset
  - ◆ 98 samples
  - ◆ 25 series from 9 popular IoT vendors
  - ◆ contains the data sets of SaTC and KARONTE

# Q1: Comparative Evaluation

## ● Overview

HemreScan raise **297** alerts with **156** vulnerabilities

## ● Effectiveness

HermeScan reports **120** more vulnerabilities than SaTC, and **152** more vulnerabilities than KARONTE

## ● Accuracy

HermeScan outperforms SaTC in TPR by **39%**

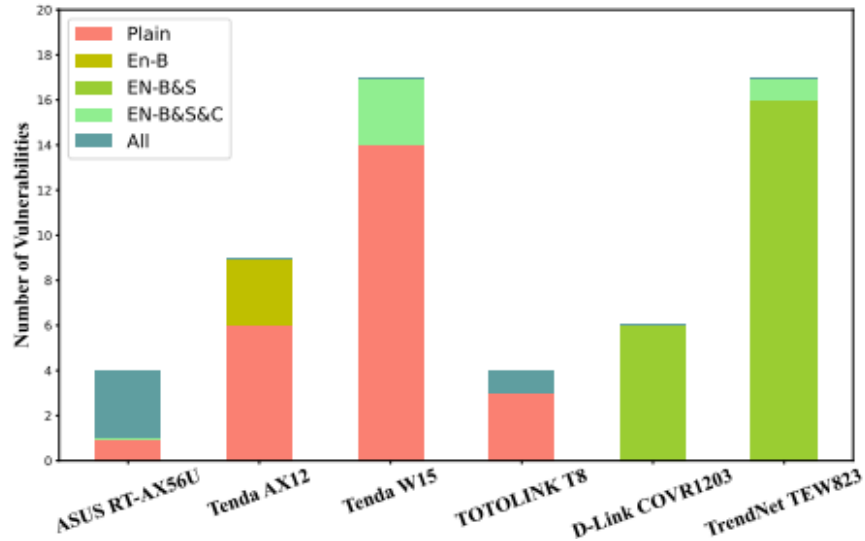
## ● Efficiency

HermeScan is **7.5x** times faster than SaTC and **3.8x** times faster than KARONTE

Vendor&Model	Program Name	HermeScan				SaTC			Karnote		
		Alerts	Vuls (bof+ci)	Vuls (other)	Time	Alerts	Vuls (bof+cli)	Time	Alerts	Vuls (bof+cli)	Time
LINKSYS MR7350	bluetoothd	0	0	0	3min	1	0	30min	0	0	2h22min
LINKSYS E9450	htpdp	0	0	0	11min	0	0	24min	0	0	2h08min
LINKSYS EA4500	twonkymediaserver	1	1	0	17min	0	0	26h	0	0	1h21min
ASUS GT-AX6000	htpdp	0	0	0	9min	5	0	27h28min	0	0	4h31min
ASUS GT-AC2900	cfg-server	0	0	0	12min	2	0	23h58min	0	0	7min
ASUS RT-AX56U	htpdp	4	4	0	13min	5	0	26h21min	0	0	48min
Tenda AX-12	htpdp	9	6	3	1h20min	0	0	12min	0	0	3h36min
Tenda AX-3	htpdp	17	11	0	2h23min	27	6	36h	0	0	1h24min
Tenda AX-1803	thttpd	20	12	2	2h03min	38	8	16h28min	0	0	5min
Tenda AX-1806	thttpd	20	14	0	2h11min	44	13	18h53min	0	0	42min
Tenda W15E	htpdp	17	15	2	2h39min	50	5	21h11min	0	0	1h22min
TOTOLINK T8	csteegi	14	4	0	14min	0	0	3min	0	0	2min
TOTOLINK LR350	csteegi	24	9	0	13min	0	0	47min	0	0	24min
TOTOLINK A7000	csteegi	18	12	0	12min	0	0	4h09min	0	0	39min
TOTOLINK A8000	csteegi	29	13	0	13min	2	0	39min	0	0	6h04min
D-LINK COVR-1201	prog.cgi	9	4	2	5h27min	0	0	10min	0	0	4h42min
D-LINK COVR-1210	prog.cgi	8	4	2	5h16min	0	0	10min	0	0	4h28min
Netgear RAX-10	net.cgi	6	0	0	19min	0	0	49min	0	0	2h20min
Netgear RAX-30	ntrg_ra_tot	0	0	0	6min	8	0	5h54min	0	0	3h10min
Netgear RAX-120	net.cgi	1	0	0	37min	0	0	1h09min	0	0	72h
Netgear MR42	htpdp	1	1	0	51min	0	0	23h56min	4	0	2h56min
Trendnet tw 829	samba_multicall	1	1	0	30min	0	0	2min	0	0	1h57min
Trendnet tew 823	ssi	39	17	0	9min	0	0	11min	0	0	1h54min
Trendnet tew 827	ssi	31	18	2	18min	0	0	27min	0	0	59min
Trendnet tew 818	rc	12	5	0	14min	0	0	18h32min	0	0	2h14min
Trendnet tew 752	cgibin	5	1	0	9min	0	0	20min	0	0	2h10min
TP-LINK AX3000	fapi_wlan_cli	0	0	0	0	0	0	15min	0	0	2h42min
TP-LINK XDR1850	dms	2	0	0	18min	0	0	3min	0	0	2min
TP-LINK XDR3060	dms	3	2	0	1h44min	0	0	2min	0	0	28min
TP-LINK XTR7880	dms	6	2	0	2h10min	0	0	3min	0	0	19min
Total	/	297	156(152)	13(11)	30h4min	182	32(32)	252h19min	4	0	127h58min
Average	/	9.9	5.2	/	1h7min	6.66	1.9	8h25min	0.13	0	4h16min

## Q2: Effectiveness of enhanced CFG

- All optimization techniques used to enhance control flow graph construction contribute to the vulnerability detection ability of HermeScan



B Function Boundary identification

S Symbol name recovery

C Shared library CFG included

# Q3: Effectiveness of Input Source Identification

- Candidate source function checking reduces the FPs of vulnerabilities by **18%** on the zero-day dataset
- Fuzzy matching strategy can find an additional **27%** of keywords
- Input source identification effectively helps HermeScan reduce false positives and false negatives

Vendor & Model	Shared Keywords(S)	Shared Keywords(H)	Increased Proportion
LINKSYS MR7350	47	52	10.64%
LINKSYS E9450	56	57	1.79%
LINKSYS EA4500	65	66	1.54%
ASUS GT-AX6000	180	187	3.89%
ASUS GT-AC2900	180	184	2.22%
ASUS RT-AX56U	404	504	24.75%
Tenda AX-12	201	222	10.45%
Tenda AX-3	246	253	2.85%
Tenda AX-1803	254	255	0.39%
Tenda AX-1806	262	269	2.67%
Tenda WISE	437	535	22.43%
TOTOLINK T8	67	69	2.99%
TOTOLINK LR350	66	67	1.52%
TOTOLINK A7000	79	80	1.27%
TOTOLINK A8000	107	116	8.41%
D-LINK COVR-1201	506	625	23.52%
D-LINK COVR-1210	495	618	24.85%
Netgear RAX-10	860	897	4.30%
Netgear RAX-30	107	237	121.50%
Netgear RAX-120	866	1005	16.05%
Netgear MR-62	862	870	0.93%
Trendnet TEW-829	35	35	0.00%
Trendnet TEW-823	1042	1459	40.02%
Trendnet TEW-827	103	365	254.37%
Trendnet TEW-818	176	249	41.48%
Trendnet TEW-752	36	36	0.00%
TP-LINK AX3000	237	238	0.42%
TP-LINK XDR1850	99	188	89.90%
TP-LINK XDR3060	96	117	21.88%
TP-LINK XTR7880	108	213	97.22%
Average	276	336	<b>27.81%</b>

# Q4: Effectiveness of Path Merging Strategy

- The path merging strategy reduces the number of paths by **89.4%** on average
- 22 out of 30 samples merged more than 90% of the paths

Vendor & Model	Paths(BF)	Paths(AF)	Decreased Proportion
LINKSYS MR7350	69	23	66.67%
LINKSYS E9450	472	121	74.36%
LINKSYS EA4500	1838	143	92.22%
ASUS GT-AX6000	1010	63	93.76%
ASUS GT-AC2900	1062	88	91.71%
ASUS RT-AX56U	636	129	79.72%
Tenda AX-12	1673	72	95.70%
Tenda AX-3	9413	96	98.98%
Tenda AX-1803	5112	109	97.87%
Tenda AX-1806	4789	101	97.89%
Tenda W15E	11113	186	98.33%
TOTOLINK T8	7126	101	98.58%
TOTOLINK LR350	12688	75	99.41%
TOTOLINK A7000	13944	77	99.45%
TOTOLINK A8000	610	105	82.79%
D-LINK COVR-1203	686	19	97.23%
D-LINK COVR-1210	644	18	97.20%
Netgear RAX-10	1299	57	95.61%
Netgear RAX-30	171	11	93.57%
Netgear RAX-120	910	345	62.09%
Netgear MR-62	2008	214	89.34%
Trendnet TEW-829	231	32	86.15%
Trendnet TEW-823	167554	265	99.84%
Trendnet TEW-827	12160	116	99.05%
Trendnet TEW-818	20014	219	98.91%
Trendnet TEW-752	594	9	98.48%
TP-LINK AX3000	0	0	0.00%
TP-LINK XDR1850	2294	17	99.26%
TP-LINK XDR3060	2747	18	99.34%
TP-LINK XTR7880	2689	18	99.33%
Average	9518	95	89.40%

# Summary

- We present a **lightweight reaching definition analysis solution** HermeScan to perform taint analysis on IoT firmware binaries
- HermeScan has discovered **87 zero-day vulnerabilities** in real-world devices, and **69** of them have been assigned **CVE IDs**
- We build **two sets of firmware samples** and comprehensively evaluate the performance of existing tools

Thanks! Questions?