



The Evolution of Program Analysis Approaches in the Era of Al

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@matrosov



- 20+ years doing all shades of binary program analysis
- Break a few times CPU's and GPU's
- Dedicating all my free time to surfing \$\mathcal{L}_\tilde{\gamma}\$

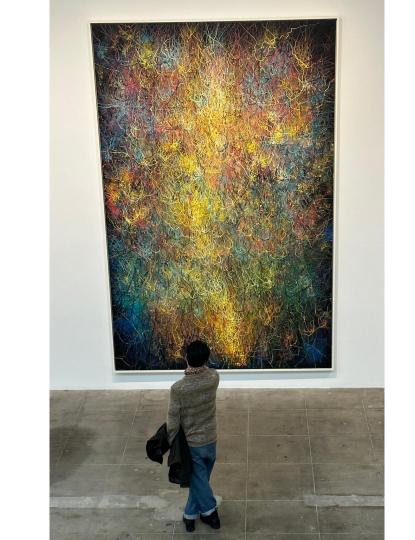


This talk is not intended to cover the complete history of binary program analysis or reverse engineering.

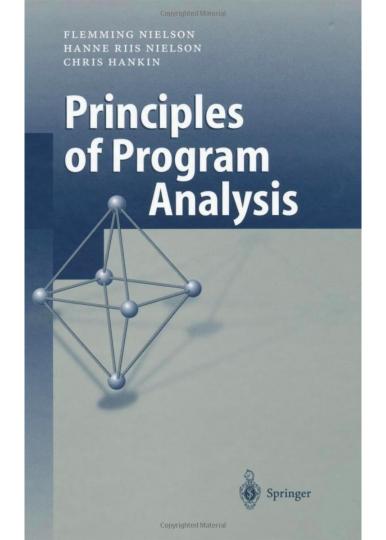


I am describing the evolution of RE from the perspective of my personal experience over 20+ years.

Academia Why is so poolly ?



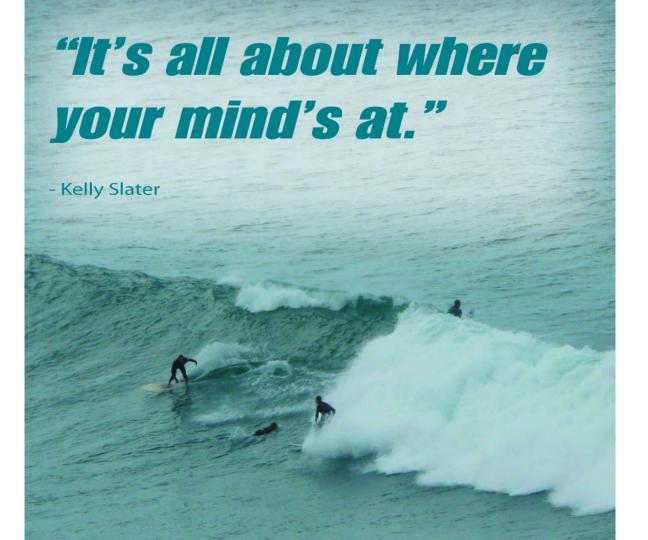




PRINCIPLES OF ABSTRACT INTERPRETATION



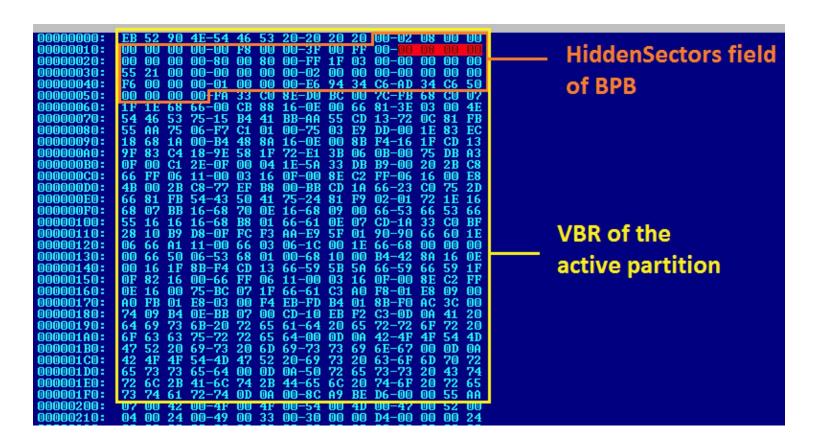
PATRICK COUSOT



major RE drivers Vulnerability research explainability coverage



Unknown structured formats

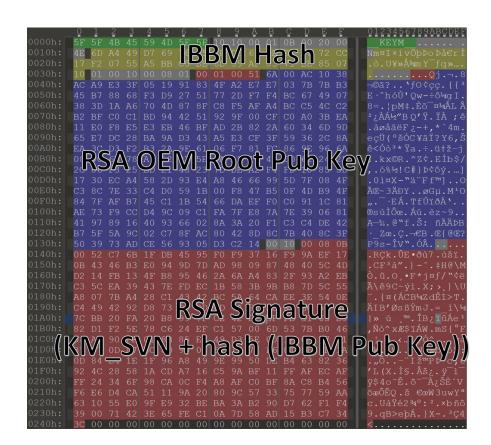


Signature-based automation

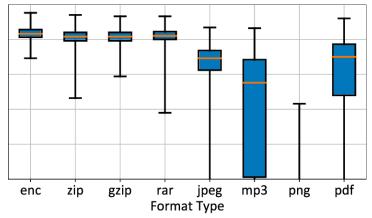


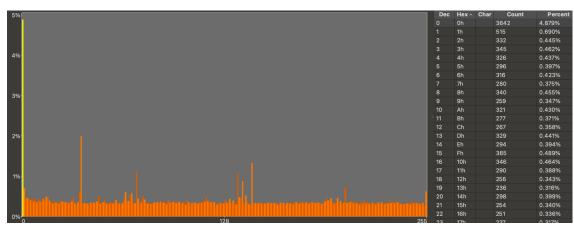
Undocumented structured formats

struct BOOT_GUARD_KEY_MANIFEST BGKM
> UBYTE Signature[8]
UBYTE Unknown
UBYTE Unknown1
UBYTE KmSvn
UBYTE Unknown2
UBYTE Unknown3
UINT16 Unknown4[0]
struct KEY_HASH IbbmKeyHash
UBYTE Unknown4[1]
UINT16 Unknown5
struct KEY_RSA OemPubKey
struct RSA_PUBLIC_KEY Key
UBYTE Unknown8
UINT16 Size
UINT32 Exp
> UBYTE PubKey[256]
UINT16 Unknown16
struct RSA_SIGNATURE Signature
UINT16 KeySize
UINT16 Unknown16
UBYTE Signature[256]



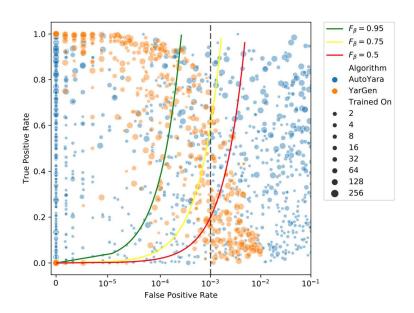
Entropy based ML models good for file format classification





Automated detection signature generation

Automatic Yara Rule Generation Using Biclustering



Signature generation Signature Optimization

alwale samp · Nossentivo. nvmbe Size of signature

Signature == sequence of bytes missing explainability no code semouties Metadoita from coole \$

Disassembly CFG => Visual Graph Representation



```
entry_point proc near
           fs, ax
            gs, ax
            eax, 4C8h
           fword ptr [eax]
           boot_guard
            ebx. eax
           edx. 0
           eax, 3
   getsec
loc_3BE6:
        dword ptr [ebp+14h], (
        eax, [ebp+8]
        short loc_3C06
            ecx, [ebp+10h]
            ecx, eax
        dl, [ecx+eax]
        dword ptr [ebp+14h]
        [eax], dl
inc
        dword ptr [ebp+14h], 0
        short loc_3BF7
     loc 3C06:
     public entry_point_1
     entry_point_1
     entry_point endp
```

REconstruction of complex C++ code still a problem in 2023

```
int thiscall Rc4 GetBufferSize( RC4 STRUCT *this)
 return (this->Reader->vTable->GetResBufSize)();
                                            int thiscall Rc4 GetBufferSize( RC4 STRUCT *this)
                                          Rc4 GetBufferSize proc near
                                                                                   ; DATA XREF:
                                                                   ecx, [ecx+4]
                                                                   eax, [ecx]
                                                                   dword ptr [eax+10h]
                                          Rc4 GetBufferSize endp
   RC4 UTABLE
                   dd offset Rc4 GetReader ; DATA XREF: sub 1011E919+1Eîo
                   dd offset Rc4 GetWriter
                   dd offset ?Destroy@EventWaitNode@details@Concurrency@@QAEXXZ
                   dd offset ?Sweep@EventWaitNode@details@Concurrency@@QAE NXZ
                   dd offset Rc4 GetBufferSize
                   dd offset Rc4 IncreaseSize
                   dd offset Rc4 Check
                   dd offset Rc4 InitEmpty
                   dd offset Rc4 Release
                   dd offset Rc4 GetMuxName
```

REIL: A platform-independent intermediate representation of disassembled code for static code analysis

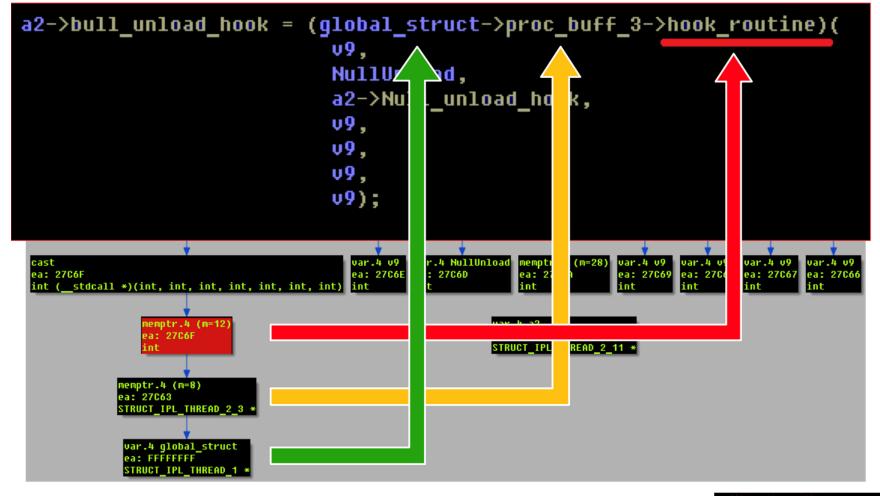
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IR => Code Semantics









HexRaysCodeXplorer v1.0: released in 2013 at REcon



```
->free mem = v4 - *v4 + 0x7D1E;
        a2->base64 encode = v4 + 0x388 - *v4;
        a2->baseó4 decode = v4 + 0x4CD - *v4:
        a2->rdtsc proc = v4 - *v4 + 0x579F:
        a2->rnd process block = v4 + 0x57A2 - *v4;
        a2->rnd fill buffer = v4 - *v4 + 0x6A87;
        a2->init rnd buffer = v4 + 0x6ABB - *v4;
        a2->field13 = v4 + 0x4B95 - *v4;
        a2->md5 init = v4 - *v4 + 0x2A2C;
  Output window
Field reference detected -> Offset 11217 : char
Field reference detected -> Offset 11218 : char
Field reference detected -> Offset 11219 : char
Field reference detected -> Offset 11220 : char
Field reference detected -> Offset 11221 : char
Field reference detected -> Offset 11222 : char
struct STRUCTURE_TYPE {
           int
                      field 0:
                      field 1;
           int
                      field 2;
           int
           int
                      field 3;
           int
                      field 4;
           int
                      field 5;
           int
                      field 6;
                      field 7;
           int
           int
                      field 8;
           int
                      field 9:
           int
                      field 10;
           int
                      field 11:
Python
```

Ghidra P-Code more suitable for RE needs vs Hex-Rays IR

```
180002aaa 42 0f b7 14 MOVZX size,word ptr [RBX + last_index*0x2]
                              R11,[0x8 + t*0x2]
180002aaf 4c 8d 1c 4d
         08 00 00 00
180002ab7 4c 03 dd
                              Decrypted, RBP
180002aba 66 83 fa 7f CMP
```

Ghidra P-Code more suitable for RE needs vs Hex-Rays IR

```
00000000180002AAA movzx edx,word ptr [rbx + r10*0x2]
unique[0x3300:8] = R10 * 0x2
unique[0x3400:8] = RBX + unique[0x3300:8]
unique [0xbe80:2] = *[ram]unique [0x3400:8]
EDX = zext(unique[0xbe80:2])
RDX = zext(EDX)
00000000180002AAF lea r11,[0x8 + rcx*0x2]
unique[0x3480:8] = RCX * 0x2
unique[0x3580:8] = 0x8 + unique[0x3480:8]
R11 = unique[0x3580:8]
00000000180002AB7 add r11,rbp
CF = carry(R11, RBP)
OF = scarry(R11, RBP)
R11 = R11 + RBP
SF = R11   < 0  
ZF = R11 == 0x0
unique[0x12e80:8] = R11 \& 0xff
unique[0x12f00:1] = popcount(unique[0x12e80:8])
unique[0x12f80:1] = unique[0x12f00:1] \& 0x1
PF = unique[0x12f80:1] == 0x0
00000000180002ABA cmp dx,0x7f
CF = DX < 0x7f
OF = sborrow(DX, 0x7f)
unique[0x28f00:2] = DX - 0x7f
SF = unique[0x28f00:2] s < 0x0
ZF = unique[0x28f00:2] == 0x0
unique[0x12e80:2] = unique[0x28f00:2] & 0xff
unique[0x12f00:1] = popcount(unique[0x12e80:2])
unique[0x12f80:1] = unique[0x12f00:1] & 0x1
PF = unique[0x12f80:1] == 0x0
```

The most comprehensive IR for RE is developed by Binary Ninja

```
000011bc EFI_STATUS sub_11bc(EFI_HANDLE* ImageHandle, struct EFI_SYSTEM_TABLE* SystemTable)
  0 @ 000011bc __saved_rsi = rsi
  1 @ 000011bc rsp = &__saved_rsi
  2 @ 000011be rsp = &var_38
  3 @ 000011c2 rax = [SystemTable + 0x58].q
  4 @ 000011c6 r8 = SignalExitBootServicesNotifier8
  5 @ 000011cd r10 = [SystemTable + 0x60].q
  6 @ 000011d1 r9 = 0
  7 @ 000011d4 [gRT1].q = rax
  9 @ 000011e0 [gRT2].q = rax
 10 @ 000011e7 rax = gSignalExitBootServicesEvent8
 11 @ 000011ee [gST].q = SystemTable
 12 @ 000011f5 SystemTable = 8
 13 @ 000011f9 [gBS1].q = r10
 14 @ 00001200 [gBS2].q = r10
 15 @ 00001207 var_18 = gSignalExitBootServicesEvent8
 16 @ 0000120c rax, ImageHandle, SystemTable, r8, r9, r10, r11, xmm4, xmm5 = call([r10 + 0x50].q, ImageHandle, SystemTable, r8, r9, stack = &var_38)
 17 @ 00001210 rax = gSignalVirtualAddressChangeEvent8
                                                                     000011bc EFI STATUS sub 11bc(EFI HANDLE* ImageHandle, struct EFI SYSTEM TABLE* SystemTable)
 18 @ 00001217 r9 = 0
 19 @ 0000121a var_18 = gSignalVirtualAddressChangeEvent8
                                                                        0 @ 000011bc __saved_rsi#1 = rsi#0
 20 @ 0000121f r8 = SignalVirtualAddressChangeNotifier8
                                                                        1 @ 000011bc rsp#1 = &__saved_rsi
                                                                        2 @ 000011be rsp#2 = &var_38
                                                                        3 @ 000011c2 rax#1 = [SystemTable#0 + 0x58].q @ mem#0
                                                                        4 @ 000011c6 r8#1 = SignalExitBootServicesNotifier8
                                                                        5 @ 000011cd r10#1 = [SystemTable#0 + 0x60].g @ mem#0
                                                                        6 @ 000011d1 r9#1 = 0
                                                                        7 @ 000011d4 [gRT1].q = rax#1 @ mem#0 -> mem#1
                                                                        9 @ 000011e0 [gRT2].q = rax#1 @ mem#1 -> mem#2
                                                                       10 @ 000011e7 rax#2 = gSignalExitBootServicesEvent8
                                                                       11 @ 000011ee [gST].q = SystemTable#0 @ mem#2 -> mem#3
                                                                       12 @ 000011f5 SystemTable#1 = 8
                                                                       13 @ 000011f9 [gBS1].q = r10#1 @ mem#3 -> mem#4
                                                                       14 @ 00001200 [gBS2].q = r10#1 @ mem#4 -> mem#5
```

15 @ 00001207 var 18#1 = gSignalExitBootServicesEvent8

Decompilation != Silver Bullet

```
void __fastcall sub_180001000(void *a1, char a2, unsigned __int64 a3)
{
   if ( a3 )
      memset(a1, a2, a3);
}
```

BinaryNinja C

Ghidra C

```
10.2.2 (9813cde2)

void FUN_180001000(undefined *param_1, undefined param_2, longlong paragrafus for (; param_3 != 0) {

for (; param_3 != 0; param_3 = param_3 + -1) {

*param_1 = param_2;
param_1 = param_1 + 1;
}

return;
}

return;
}
```

The decompilation of Golang is a disaster

```
__int64 v22; // x28
<u>__int64 v23; // x0</u>
__int64 v24; // x1
unsigned __int64 v26; // x0
__int64 v27; // x3
__int64 v28; // x4
__int64 v29; // x5
__int64 v30; // x6
__int64 v31; // x7
__int64 v33; // [xsp+8h] [xbp-60h]
__int64 v34; // [xsp+10h] [xbp-58h]
__int64 v35; // [xsp+18h] [xbp-50h]
int64 v36; // [xsp+20h] [xbp-48h]
unsigned __int64 v37; // [xsp+50h] [xbp-18h]
int64 v38; // [xsp+58h] [xbp-10h]
while ( \&a9 <= *(v22 + 16) )
 a10 = a1:
 runtime_morestack_noctxt_abi0();
 a1 = a10
 a2 = a11
a10 = a1
v34 = strings_Replace();
v38 = runtime_stringtoslicebyte(0LL, v23, v24);
v26 = encoding_hex_Decode();
if (v37 < v26)
 runtime_panicSliceAcap();
return bytesToGUID(v38, v26, v37, v27, v28, v29, v30, v31, v33, v34, v35, v36);
```

The decompilation of Rust is a disaster

```
v7 = atomic_load(&uefi::parsers::bg::KEYM_TAG::he79a90e6d0585aaf);
1f ( v7 != 2 )
  once_cell::imp::OnceCell$LT$T$GT$::initialize::h0ec87ed7e59f3858(
    &uefi::parsers::bg::KEYM_TAG::he79a90e6d0585aaf,
    &uefi::parsers::bg::KEYM_TAG::he79a90e6d0585aaf);
regex::re_bytes::Regex::find_iter::hef832e90fc36d1e1(&unk_10019C008, a1, a2);
v45 = v41;
v46 = v42;
v47 = v43:
v48 = v44:
_$LT$alloc..vec..Vec$LT$T$GT$$u20$as$u20$alloc..vec..spec_from_iter..SpecFromIter$LT$T$C$I$GT$$GT$::from_iter::h28d513b9e8021fff(&v45);
result = alloc::slice::merge_sort::h80d655c35f129dbe();
if ( v40 >= 2 )
  v31 = v40 - 1;
  v32 = (unsigned __int64 *)(v38 + 8);
  V8 = 1LL:
  do
    v10 = +v32;
    if (*v32 != *(_QWORD *)(v38 + 8 * v8 - 8))
      *(QWORD *)(V38 + 8 * V8++) = V10;
    ++v32:
  while ( V31 ):
if ( IV8 )
ABEL 42:
  if ( v39 )
    result = __rust_dealloc(v38, 8 * v39);
  a3[2] = OLL;
  a3[3] = 0LL;
  *a3 = 11LL:
  a3[1] = 1LL;
  return result;
v11 = 8LL:
v12 = 8 + v8;
v13 = 1LL;
v14 = 14LL;
v36 = 14LL:
v37 = 1LL;
while (1)
  result = *(_QWORD *)(v38 + v11);
  if ( result > a2 )
    core::slice::index::slice_start_index_len_fail::h2a4533b191a8042c();
```



Next directions for REsearch

 Utilize more Data Flow Analysis, industry is too focused on Control Flow Analysis and missing out on Data Semantics.

Data and Code Reconstruction required specific methods to preserve code and its dependencies, as well as fast methods of querying this data. Datalog can be used to represent data and code in a deductive database, but it requires a large amount of memory.

 Infer ML models based on code semantics, not byte sequences, which lack context.



{* SECURITY *}

One month after Black Hat disclosure, HP's enterprise kit still unpatched

What could go wrong with leaving firmware open after world's biggest hacker convention talk?

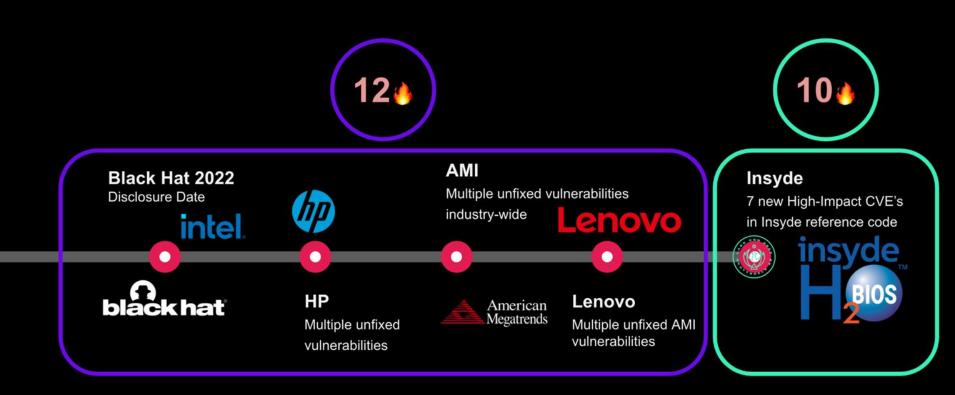


Firmware Repeatable Failures



Vendor	Vulnerabilities	Number of Issues	BINARLY ID	CVE ID	CVSS score
intel.	PEI Memory Corruption (Arbitrary Code Execution)	3	BRLY-2022-027 BRLY-2022-009 BRLY-2022-014	CVE-2022-28858 CVE-2022-36372 CVE-2022-32579	8.2 High 8.2 High 7.2 High
<u></u> ami	DXE Arbitrary Code Execution	1	BRLY-2022-015	CVE-2022-34345	7.2 High
<u> </u>	SMM Memory Corruption (Arbitrary Code Execution)	2	BRLY-2022-003 BRLY-2022-016	CVE-2022-27493 CVE-2022-33209	7.5 High 8.2 High
(hp	SMM Memory Corruption (Arbitrary Code Execution)	6	BRLY-2022-010 BRLY-2022-011 BRLY-2022-012 BRLY-2022-013 BRLY-2021-046 BRLY-2021-047	CVE-2022-23930 CVE-2022-31644 CVE-2022-31645 CVE-2022-31646 CVE-2022-31640 CVE-2022-31641	8.2 High 7.5 High 8.2 High 8.2 High 7.5 High 7.5 High

Firmware Security Repeatable Failures



Vulnerabilities in the Insyde (industry-wide impact)

BRLY	CVE	CVSS v3	Description
BRLY-2022-017	CVE-2022-36338	7.5 High	SMM callout vulnerability in SMM driver (SMM arbitrary code execution)
BRLY-2022-018	CVE-2022-35894	6.0 Medium	SMM memory leak vulnerability in SMM driver (SMRAM read)
BRLY-2022-019	CVE-2022-36337	7.7 High	The stack buffer overflow vulnerability in DXE driver
BRLY-2022-020	CVE-2022-35407	7.7 High	The stack buffer overflow vulnerability in DXE driver EKOPARTY
BRLY-2022-021	CVE-2022-35897	7.7 High	The stack buffer overflow vulnerability in DXE driver
BRLY-2022-022	CVE-2022-35408	7.5 High	SMM callout vulnerability in SMM driver (SMM arbitrary code execution)
BRLY-2022-023	CVE-2022-36448	8.2 High	SMM memory corruption vulnerability in Software SMI handler
BRLY-2022-024	CVE-2022-35895	8.2 High	SMM memory corruption vulnerability in SMM driver (SMRAM write)
BRLY-2022-025	CVE-2022-35896	6.0 Medium	SMM memory leak vulnerability in SMM driver (SMRAM read)
BRLY-2022-026	CVE-2022-35893	8.2 High	SMM memory corruption vulnerability in SMM driver (SMRAM write)



- The REsearch year in numbers:
- 🖐 Total number of vulnerabilities reported 228 🤲
- XAffected silicon vendors Intel, AMD, Qualcomm
- 💥 Affected IBVs Insyde, AMI
- *Affected device vendors MS, HP, HPE, Dell, Lenovo, Intel, Fujitsu, Framework, Atos, Aruba, Cisco, Juniper ...

3:07 PM · Dec 28, 2022 · **33.9K** Views



COMPCEXIL



Revisiting Automated Bug Hunting

Progression of our past work:

"efiXplorer: Hunting for UEFI Firmware Vulnerabilities at Scale with Automated Static Analysis" 1

- Scalable approach based on vulnerability models; combination of:
 - 1. Lightweight static analysis
 - 2. Under-constrained symbolic execution

Limitations of current approaches

With great scalability, comes a (great) potential for false positives!

```
Address Type

0000000FFAE2BFD pei_get_variable_buffer_overflow

0000000FFAE8894 pei_get_variable_buffer_overflow
```

```
push
push
        ecx, offset EFI PEI READ ONLY VARIABLE2 PPI GUID
call
        eax, [ebp+Data]
                        : Data
                          DataSize
                        ; Attributes
        offset EFI_SETUP_VARIABLE_GUID ; VariableGuid
push
        offset VariableName : "Setup'
push
        [eax+EFI_PEI_READ_ONLY_VARIABLE2_PPI.GetVariable] ; VariablePPI->GetVariable()
                        ; EFI STATUS (EFIAPI *EFI PEI GET VARIABLE2)(IN CONSTEFI PEI READ ONLY VARIABLE2 PPI
                        ; Attributes
push
        offset stru_FFAEE1D0 ; VariableGuid
        offset aPchsetup : "PchSetup"
push
        [eax+EFI_PEI_READ_ONLY_VARIABLE2_PPI.GetVariable] ; VariablePPI->GetVariable()
                          EFI STATUS (EFIAPI *EFI PEI GET VARIABLE2)(IN CONSTEFI PEI READ ONLY VARIABLE2 PPI
```

```
push
                         : Ppi
                         : PpiDescriptor
push
                         : Instance
push
        offset EFI_PEI_READ_ONLY_VARIABLE2_PPI_GUID ; Guid
push
                         : PeiServices
call
        [eax+EFI_PEI_SERVICES.LocatePpi] ; qPS->LocatePpi()
                         ; EFI_STATUS(EFIAPI * EFI_PEI_LOCATE_PPI) (IN CONST EFI_PEI
        [ebp+DataSize],
lea
        eax, [ebp+Data]
                         ; Data
lea
push
                         : DataSize
push
                         : Attributes
push
        offset stru_FFAEE230 ; VariableGuid
        offset aSasetup : "SaSetup"
push
call
        [eax+EFI_PEI_READ_ONLY_VARIABLE2_PPI.GetVariable] ; VariablePPI->GetVariable
                          EFI STATUS (EFIAPI *EFI PEI GET VARIABLE2)(IN CONSTEFI PEI
lea
push
lea
push
                          DataSize
mov
                         : Attributes
push
        offset stru_FFAEDF70 ; VariableGuid
push
        offset aCpusetup
push
                         ; This
call
        [eax+EFI_PEI_READ_ONLY_VARIABLE2_PPI.GetVariable] ; VariablePPI->GetVariable
                        ; EFI_STATUS (EFIAPI *EFI_PEI_GET_VARIABLE2)(IN CONSTEFI_PEI
        eax, [ebp+var_688]
lea
mov
        ecx esi
push
```

Limitations of current approaches

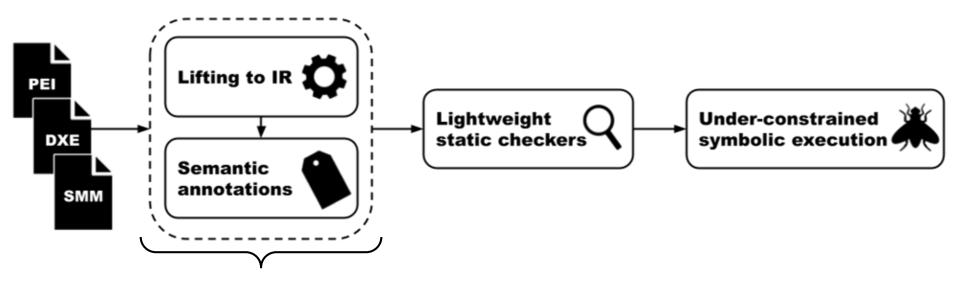
Limitations of existing approaches:

- Large number of false positives
- Mostly based on syntactic properties (pattern matching on disassembly)
- Highlighted in research by SentinelOne (Brick²):
 - Pattern matching on decompiler output
 - But: requires decompiler (Hex-Rays) & will not scale

Binarly team approach:

- Leverage semantic properties
- Use lightweight code pattern *checkers* to provide hints for deeper analysis

Analysis pipeline

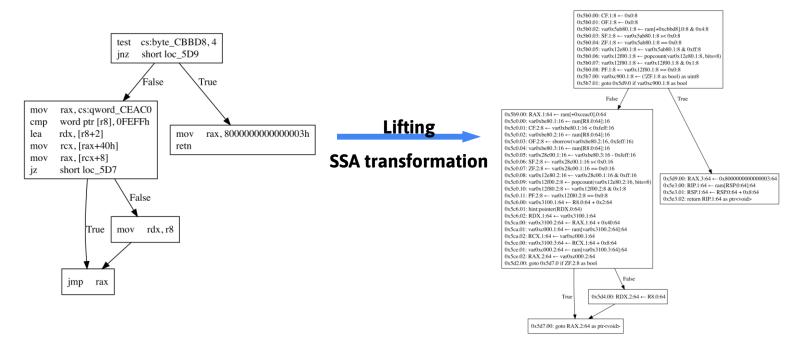


Typically takes 4-6s per firmware image (100s of modules)

Inspired by: "Sys: A Static/Symbolic Tool for Finding Good Bugs in Good (Browser) Code" (Brown et al., USENIX Security 2020)

IR lifting

- Extract uniform SSA from IR representation for 32-bit and 64-bit modules
- IR explicitly encodes instruction side-effects



Binarly Semantic annotations

```
•
```

```
binarly::efi::services
                        service call to InstallPpi: EFI_PEI_INSTALL_PPI
                        resolved type: ptr<fn(PeiServices: ptr<PEFI_PEI_SERVICES>, PpiList: ptr<EFI_PEI_PPI_DESCRIPTOR>) -> EFI_STATUS>
binarly::efi::services]
binarly::efi::services] - PeiServices: ptr<PEFI_PEI_SERVICES> = 0xfadefada:32
binarly::efi::services] - PpiList[0]: struct<EFI PEI PPI DESCRIPTOR>
binarly::efi::services]
                          - Flags: 0x10:32
binarly::efi::services]
                          - Guid: EFI_PEI_RESET_PPI_GUID
binarly::efi::services]
                                   0xffac4a3c
                          - Ppi:
                        - PpiList[1]: struct<EFI_PEI_PPI_DESCRIPTOR>
binarly::efi::services]
binarly::efi::services]
                          - Flags: 0x80000010:32
                                  AMI_PEI_SBINIT_POLICY_PPI_GUID
binarly::efi::services]
                          - Guid:
binarly::efi::services]
                                   0xffac4a38
                          - Ppi:
```

- Annotate IR with types and service information (similar to efiXplorer³ and FwHunt⁴)
- Identify analysis entry-points based on module type, e.g.:
 - SMI handlers (DXE/SMM modules)
 - PEI notification callbacks (PEI modules)

^{3: &}lt;a href="https://github.com/binarly-io/efiXplorer">https://github.com/binarly-io/efiXplorer

^{4:} https://github.com/binarly-io/fwhunt-scan

Binarly Static checkers

Checkers based on lightweight static analysis defined using an eDSL:

```
let mut matcher_builder = MatcherBuilder::new();
let s1 = matcher_builder.add_rule(ServiceCall::new(&project, "GetVariable"));
let s2 = matcher_builder.add_rule(ServiceCallChain::new(&project, "GetVariable"));
matcher_builder.add_transition(s1, s2)?;
matcher_builder.add_terminal(s2);
```

- Control-flow properties (reachability)
- Data-flow properties (data-dependence)
- Inferred call-site properties (e.g., arguments passed, type information)
- Domain-specific annotations:
 - Service-specific (e.g., GetVariable variants in PEI and DXE phases)
 - Common APIs (e.g., CopyMem, ZeroMem, etc.)

Under-constrained Symbolic Execution

Similar to past research:
 "Finding BIOS Vulnerabilities with Symbolic Execution and Virtual Platforms"

Binarly team approach:

- Instrument anything (IR operation granularity)
- Simulate execution from anywhere
- Reason about hardware interactions and partial state using symbolic variables injected during simulation
- Identify violations of model assumptions (e.g., input to API should not be user-controlled)
- No source-code required!

PEI-phase vulnerabilities



```
target/release/peiscan -v -d data -e EFI_PEI_END_0F_PEI_PHASE_PPI_GUID ./SbPei-c1fbd624-27ea-40d1-aa48-94c3d.
```

(BRLY-2022-014/CVE-2022-32579)

GetVariable leading to arbitrary write

PEI-phase vulnerabilities

./target/release/peiscan -v -d data PlatformInitAdvancedPreMem-56bbc314-b442-4d5a-ba5c-d842dafdbb24.peim

(BRLY-2022-027/CVE-2022-

binarly checkers::types reached terminal for this path

(base)

```
28858)
                     GetVariable without DataSize check
                              False Positive detection
binarly_checkers::types] setting label for rule 0 on entity 54 at 0xffae8894
binarly_checkers::types] setting label for rule 0 on entity 157 at 0xffae8871
binarly_checkers::types] setting label for rule 1 on entity 54 at 0xffae8894
binarly_checkers::types] setting label for rule 1 on entity 157 at 0xffae8871
binarly_checkers::types] stepping the searcher
binarly_checkers::types] no current checker
binarly_checkers::types] new checker has length 2
binarly_checkers::types] rule state 0 matches entity 54
binarly_checkers::types] rule state 0 accepts entity 54
binarly_checkers::types] continue with next transition
binarly_checkers::types] rule state 1 matches entity 54
binarly_checkers::types] rule state 1 does not accept transition to entity 54
binarly_checkers::types] rule state 1 matches entity 157
binarly_checkers::types] rule state 1 does not accept transition to entity 157
binarly_checkers::types] removing last transition set
binarly_checkers::types] rule state 0 matches entity 157
binarly_checkers::types] rule state 0 accepts entity 157
binarly_checkers::types] continue with next transition
binarly checkers::types rule state 1 matches entity 54
binarly_checkers::types] rule state 1 accepts entity 54
```

DXE/SMM vulnerabilities



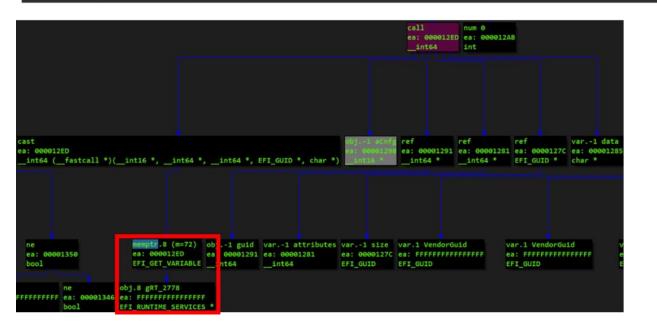
(BRLY-2022-016/CVE-2022-33209)

```
./target/release/smiscan -v -d data ./SmmSmbiosElog-8e61fd6b-7a8b-404f-b83f-aa90a47cabdf.smm
```

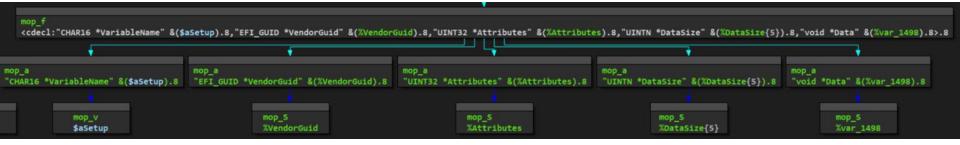
Buffer overflow discovery &

CommBuffer reconstruction

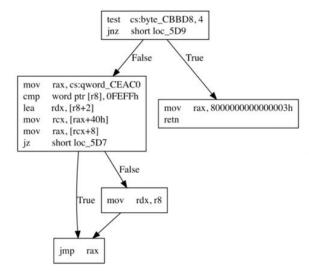
[gRT_2778->GetVariable)(aCnfg, &guid, &attributes, &size, data)







Disassembly



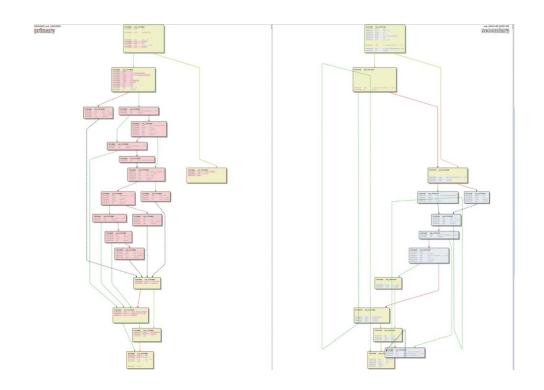
Lifting & SSA Transformation



Embedding

0.095785	-0.015778	-0.079486
0.059728	0.028905	0.01277
-0.044367	-0.052569	0.011392
-0.0086491	0.02391	-0.050848
-0.013871	0.0006036	7 0.02299
-0.054943	0.066296	-0.019087
-0.062606	0.14307	0.0084581
-0.01847	0.038296	-0.061336
-0.079965	-0.042986	-0.027591
0.095317	0.045197	0.099199
0.040439	-0.080677	-0.00061382
0.089344	-0.076245	0.052956
-0.019518	-0.064788	-0.059764
-0.03483	-0.051194	0.0042634
-0.033321	0.028235	0.031004
0.049709	-0.037423	0.024112
0.068241	0.043215	0.099272
0.13301	-0.038987	0.051024
0.065909	-0.020939	0.051219
-0.050137	-0.040482	0.035888
-0.015513	-0.044076	0.044773
-0.051152	0.049211	0.0056971
0.026995	0.064005	0.025534
-0.03215	-0.11745	0.01306
-0.045706	0.0091048	-0.019097
0.044011	0.0043315	-0.021892
-0.080179	-0.045489	-0.016057
0.063371	-0.11101	0.066997
-0.012043	0.020092	0.032347
-0.0059101	0.032843	0.047494
0.0024613	0.022228	0.022552
-0.072352	0.020193	-0.024909
0.062153 -	0.016538	0.0045914

Binary Diffing == BinDiff



"Graph-based comparison of Executable Objects" - 2005, SSTIC

https://actes.sstic.org/SSTIC05/Analyse_differentielle_de_binaires/SSTIC05-article-Flake-Graph_based_comparison_of_Executable_Objects.pdf

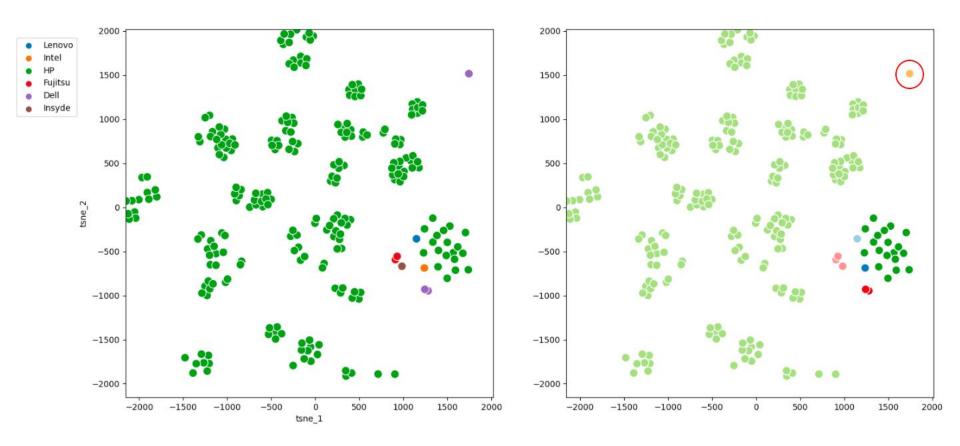
Binary Diffing == BinDiff

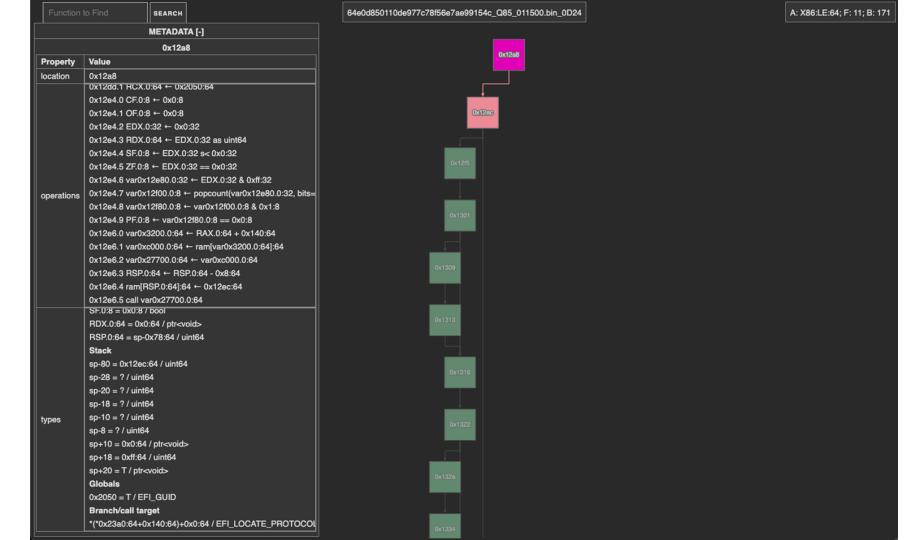
• • •	Matched Functions					
Similarity	Confid	Change	EA Primary	Name Primary	EA Secondary	Name Secondary
0.95	0.99	GI	000000018001	CoreCreateEventInternal	000000018001F10C	sub_000000018001F10C
0.99	0.99	-I	000000018000	CoreExitBootServices	0000000180009F4C	sub_0000000180009F4C
0.99	0.99	-I	000000018002	InternalAllocatePool	00000001800233B0	sub_00000001800233B0
1.00	0.99		000000018000	sub_18000030C	000000018000030C	sub_000000018000030C
1.00	0.99		000000018000	sub_180000358	0000000180000358	sub_0000000180000358
1.00	0.99		000000018000	sub_180000420	0000000180000420	sub_0000000180000420
1.00	0.99		000000018000	sub_18000045C	000000018000045C	sub_000000018000045C
1.00	0.99		000000018000	sub_1800004EC	00000001800004EC	sub_00000001800004EC
1.00	0.99		000000018000	sub_1800005D0	00000001800005D0	sub_00000001800005D0
1.00	0.99		000000018000	sub_180000624	0000000180000624	sub_0000000180000624
1.00	0.99		000000018000	sub_1800006D8	00000001800006D8	sub_00000001800006D8
1.00	0.99		000000018000	sub_180000730	0000000180000730	sub_0000000180000730
1.00	0.99		000000018000	sub_1800007A4	00000001800007A4	sub_00000001800007A4
1.00	0.99		000000018000	sub_1800007E4	00000001800007E4	sub_00000001800007E4
1.00	0.99		000000018000	sub_1800007FC	00000001800007FC	sub_00000001800007FC
1.00	0.99		000000018000	sub_180000910	0000000180000910	sub_0000000180000910
1.00	0.99		000000018000	sub_1800009E0	00000001800009E0	sub_00000001800009E0
1.00	0.99		000000018000	sub_180000A4C	0000000180000A4C	sub_0000000180000A4C
1.00	0.99		000000018000	sub_180000AD4	0000000180000AD4	sub_0000000180000AD4

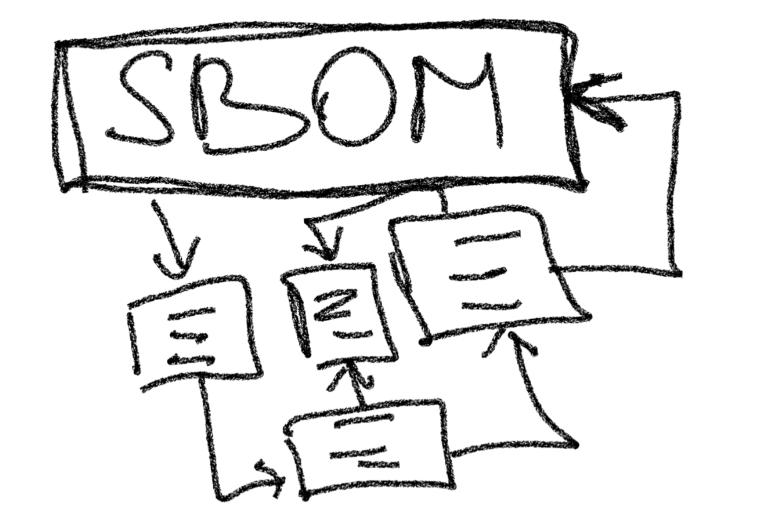
```
"name": "FunctionSimilarity",
"<u>meta</u>": {
  "description": "Check how similar the module's functions are to the same module (by GUID) in a previous firmware version.",
  "extra info": {
    "modules": [
        "guid": "d6a2cb7f-6a18-4e2f-b43b-9920a733700a",
        "hash": "9df301ebb3d4035ff173a0f66c17f1fa8c01241b7a472b9fce5927b1019c9eed",
        "name": "DxeCore",
        "<u>similarity</u>": "Very dissimilar (less than 20% similarity)"
  "<u>severity</u>": 2
"status": 1
"name": "AddedModuleVariables",
"meta": {
  "<u>description</u>": "Check if this module references any variables that were not referenced by the same module (by GUID) in a previous firmware version.",
  "extra info": {
    "modules": [
        "guid": "d6a2cb7f-6a18-4e2f-b43b-9920a733700a",
        "name": "DxeCore",
        "new variables": [
```



T-distributed Stochastic Neighbor Embedding (TSNE)

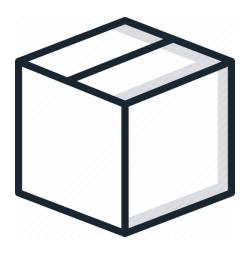




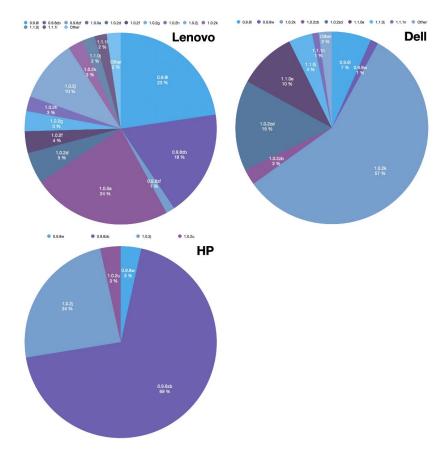


SBOM => Open/Closed Source Challenges





SBOM == Policy != Technology



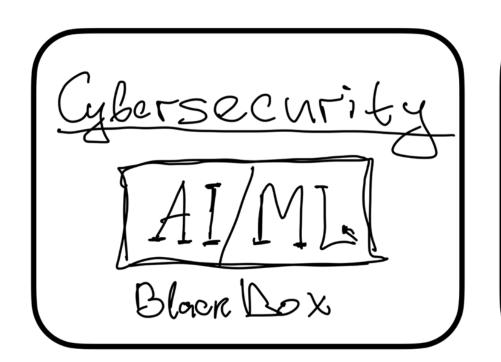
Next directions for REsearch

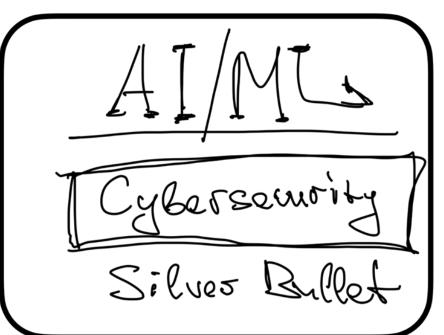
Detecting known vulnerabilities is different from finding known unknowns.
 When automating vulnerability research, it is extremely important to scope the search area correctly.

We find more problems than we can automatically explain and triage.
 Automating the process of explaining exploitability of the findings is one of the most important challenges facing the industry.

 ML models guided by code semantics can automate the search for well documented security problems.

The new old challenges of machine learning





AI/ML doesn't solve all problems magically







Thank You!