







Predictive Context-sensitive Fuzzing

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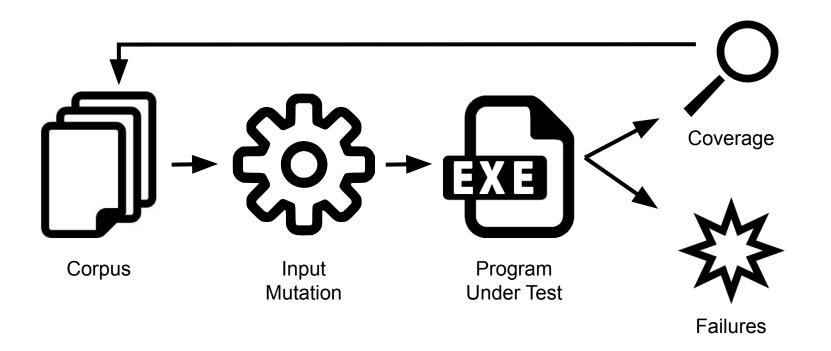








Coverage-guided fuzzing







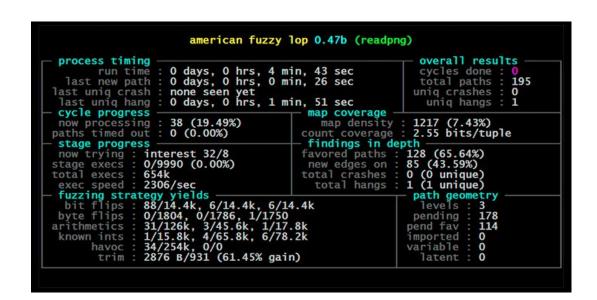


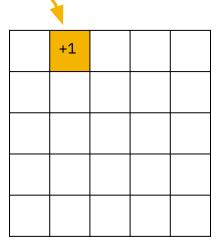


Edge-coverage guided fuzzing

hash_edge(0x41414141, 0x42424242)

0x41414141: jmp 0x42424242





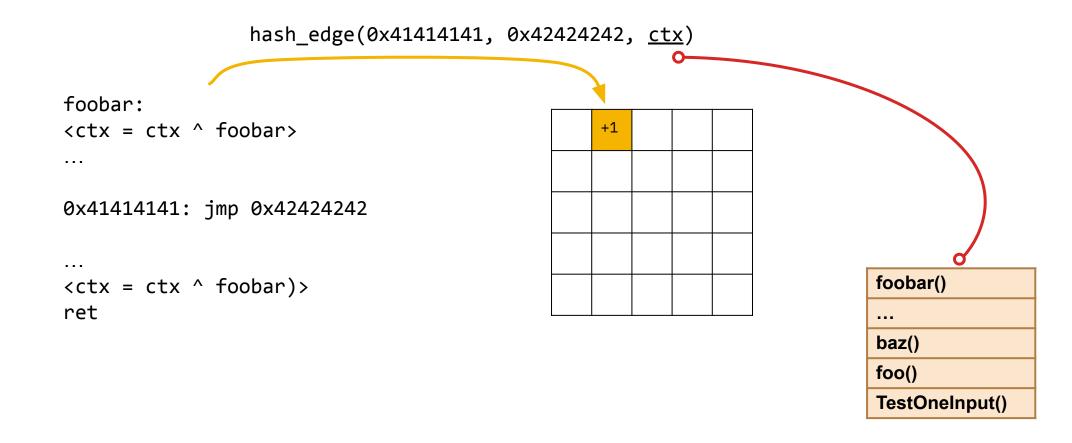








Context-sensitive, edge-coverage guided fuzzing











Why helpful?

```
void cmd_handler_foo(int a, size_t b) { memset(buf, a, b); }
void cmd handler bar(int a, size t b) { cmd handler foo(...
void cmd handler baz(int a, size t b) { cmd handler bar(...
typedef void (* dispatch t)(int, size t);
dispatch_t handlers[UCHAR_MAX] = {
   cmd_handler_foo,
   cmd_handler_bar,
   cmd handler baz,
};
int main(int argc, char **argv)
   int cmd;
   while ((cmd = getchar()) != EOF) {
       if (handlers[cmd]) {
            handlers[cmd](getchar(), getchar());
```

coverage of cmd bar is a superset of cmd foo

Project Zero

News and updates from the Project Zero team at Google

Wednesday, December 1, 2021

This shouldn't have happened: A vulnerability postmortem

Posted by Tavis Ormandy, Project Zero

Introduction

This is an unusual blog post. I normally write posts to highlight some hidden attack surface or interesting complex vulnerability class. This time, I want to talk about a vulnerability that is neither of those things. The striking thing about this vulnerability is just how simple it is. This should have been caught earlier, and I want to explore why that didn't happen.

In 2021, all good bugs need a catchy name, so I'm calling this one "BigSig".

First, let's take a look at the bug, I'll explain how I found it and then try to understand why we missed it for so long.









Challenges to efficiency

- 1 coverage map explosion
- queue explosion

Fuzzer	Queue size	Exec/sec	Collisions (map use)
classic edge (2^16 map)	9 911	609	9.8% (19.86%)
pcguard edge (collfree)	11 093	572	-
ctx-sensitive (2^16 map)	33 675	530	50.7% (79.54%)
ctx-sensitive (2^20 map)	21 157	1 84	1.2% (7.21%)
predictive ctx-sensitive	15 455	490	-

24h on **libxml2** - AFL++ 3.15a Drivers & seeds: FuzzBench







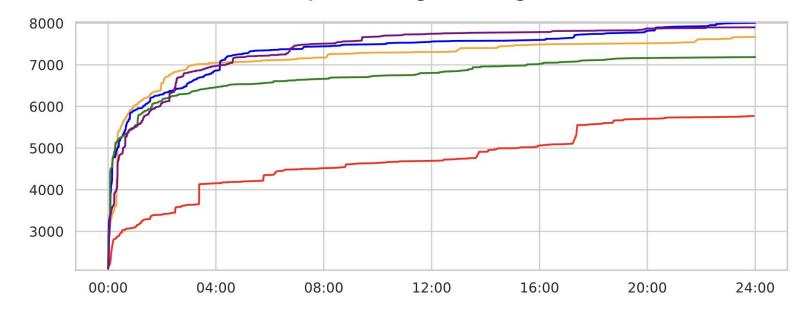


Challenges to efficiency

- coverage map explosion
- queue explosion

Fuzzer classic edge (2^16 map) pcguard edge (coll.-free) ctx-sensitive (2^16 map) ctx-sensitive (2^20 map) predictive ctx-sensitive

Impact on edge coverage











Making it efficient...

Key ideas

- 1. encode context without collisions
- 2. track context only at selected regions
- 3. predict profitable regions

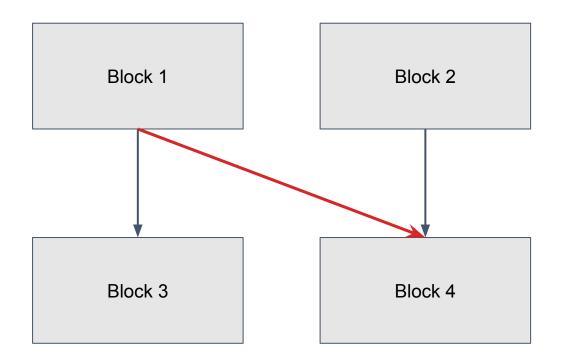








Collision-free edge coverage



Splitting CFG critical edges removes collisions in edge-coverage tracking

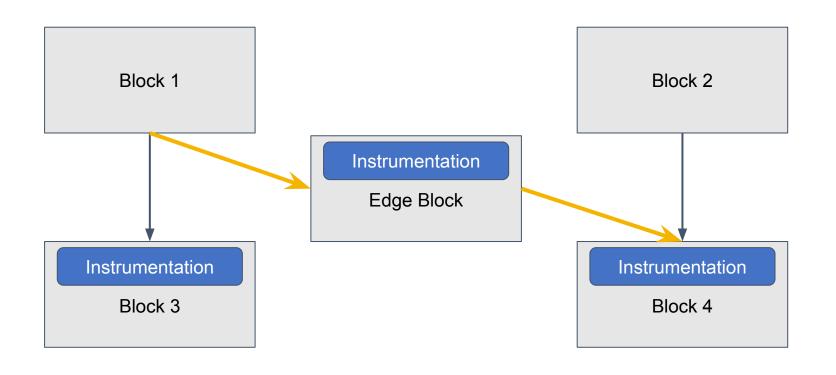








Collision-free edge coverage



As basic-block IDs now suffice to uniquely identify edges

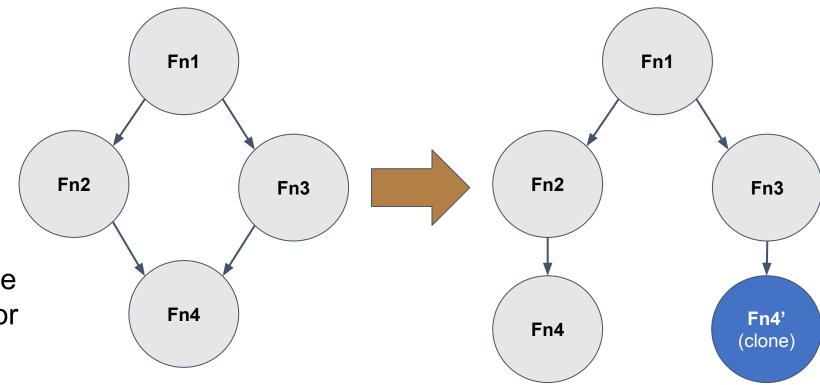








Collision-free context-sensitivity?



By **cloning** a function, we get unique CFG edges for a (caller, callee) pair









Selective sensitivity

Cloning is fuzzer-friendly, but the path explosion problem stays!

Benchmark (FuzzBench)	Edges	Functions	Call sites	Calling contexts
ffmpeg	716 K	5 K	44 K	8 M
libarchive	67 K	<1 K	4 K	27 M
libhevc	120 K	2k	< 1K	125 M
libxml2	104 K	1k	7 K	44 B
njs	57 K	490	4 K	13 M









Initial strategies

Prioritize within cloning budget

- favor call sites from nodes closer to call-graph root (harness)
- favor call sites closer to leaves
- treat every call site with the same priority

vs. random selection











Data flow-based prediction

Prioritize call sites with distinctive incoming data

- clone one if it passes data "seen less often" at other callers
- diversity as a proxy for interesting
- needs only call-site sensitivity (i.e., no full contexts)
- focus on pointer-type arguments

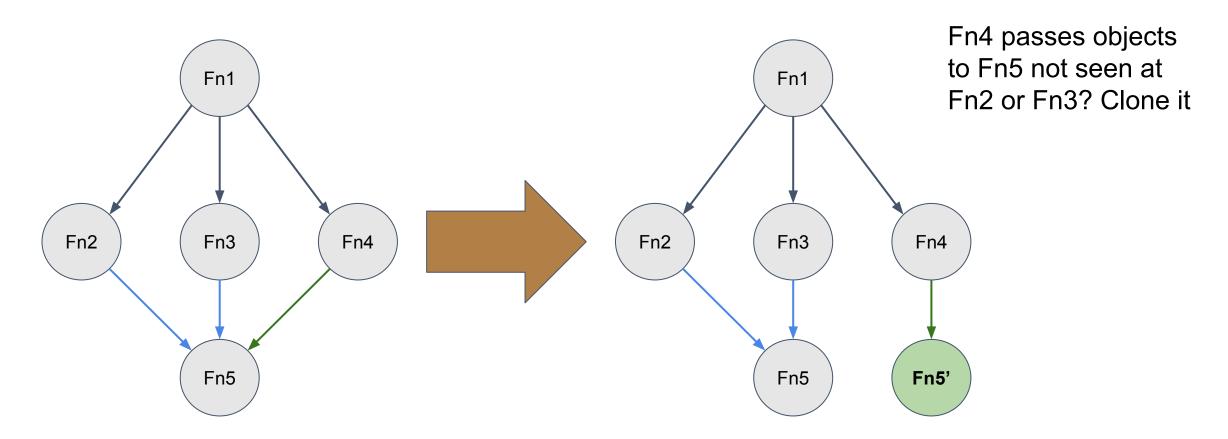








Data flow-based prediction











Implementation

- gllvm
- AFL++ 3.15a
- LLVM 10
- SVF framework

https://github.com/eurecom-s3/predictive-cs-fuzzing













Evaluation

RQ: Can we find bugs that existing approaches overlook?

RQ: How is fuzzing performance affected?

Who: pcguard LTO, Angora-style CS fuzzing, Predictive CS fuzzing

How: 16 programs from FuzzBench bug benchmarks

budget of 2¹⁸ entries (L2 size)









Bug counts

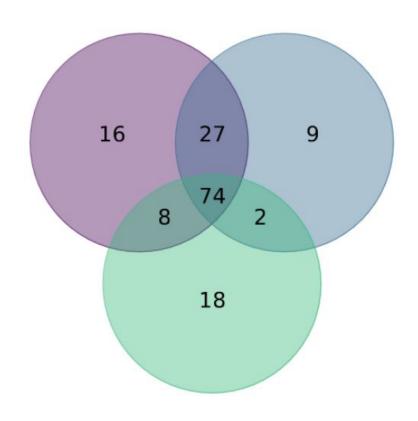
predictive (125)

Ito (112)

context (102)

Highlights

- +11.6% than Ito, +22.5% than context
- 23 of our bugs (19.2%) were **missed** by Ito (7 from new coverage, 16 from exploitation)











And new bugs!

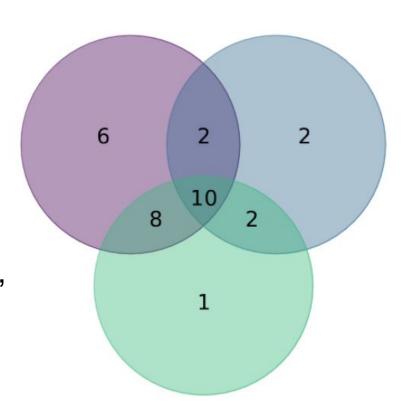
predictive (26)

Ito (16)

context (21)

Highlights

- 26 out of 31 were exposed by predictive
- 8 security issues (1 each in ffmpeg, njs, libhevc, and matio; 4 in stb) 6 CVEs assigned











Fuzzing performance

Trends

- queue size: +26.4% vs lto (context: +81.7%)
- throughput: 6.5% slower than Ito (context: 20.3%)
- coverage: close to Ito on 12/16 subjects, better on 8
- tenable compilation costs, 3.6x binary size increase









Final remarks

Existing approaches face an **impossible trade-off** between collisions and trashing from queue/map explosion. We show a profitable avenue as we proactively select the most promising contexts based on data-flow diversity

Future opportunities: non-pointer arguments, cloning for indirect calls